

Static Program Analysis Meets Test Case Generation

Maria Christakis
MPI-SWS, Germany

About me

About me



1986

Born in Crete,
Greece

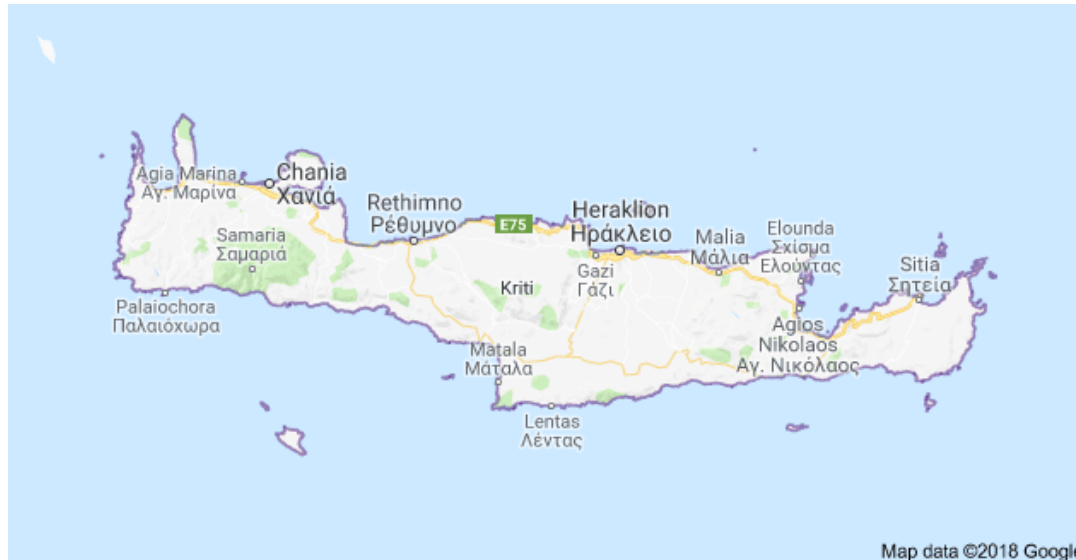
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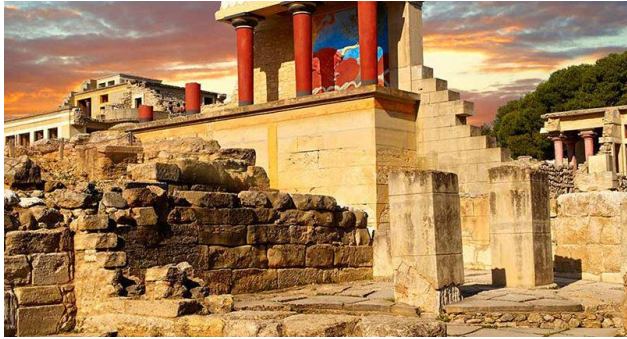
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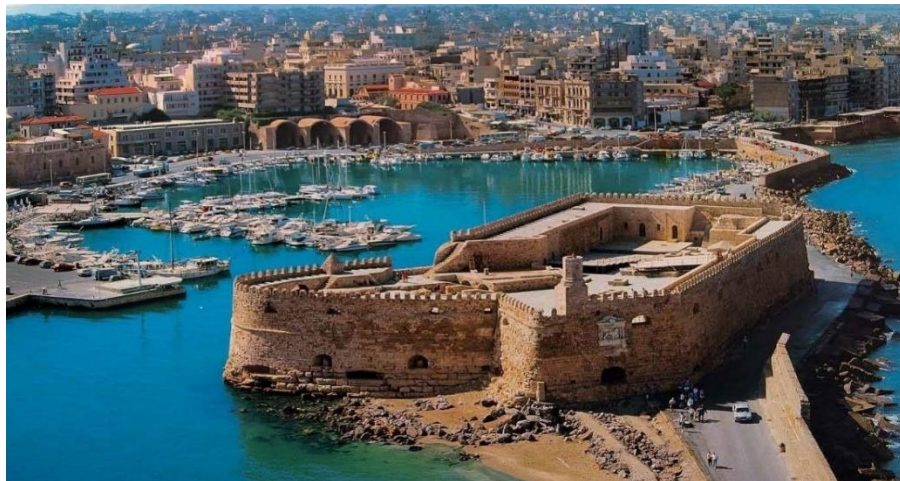
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2003 – 2009

Dipl. in Electrical and Computer Engineering
N.T.U.A., Greece



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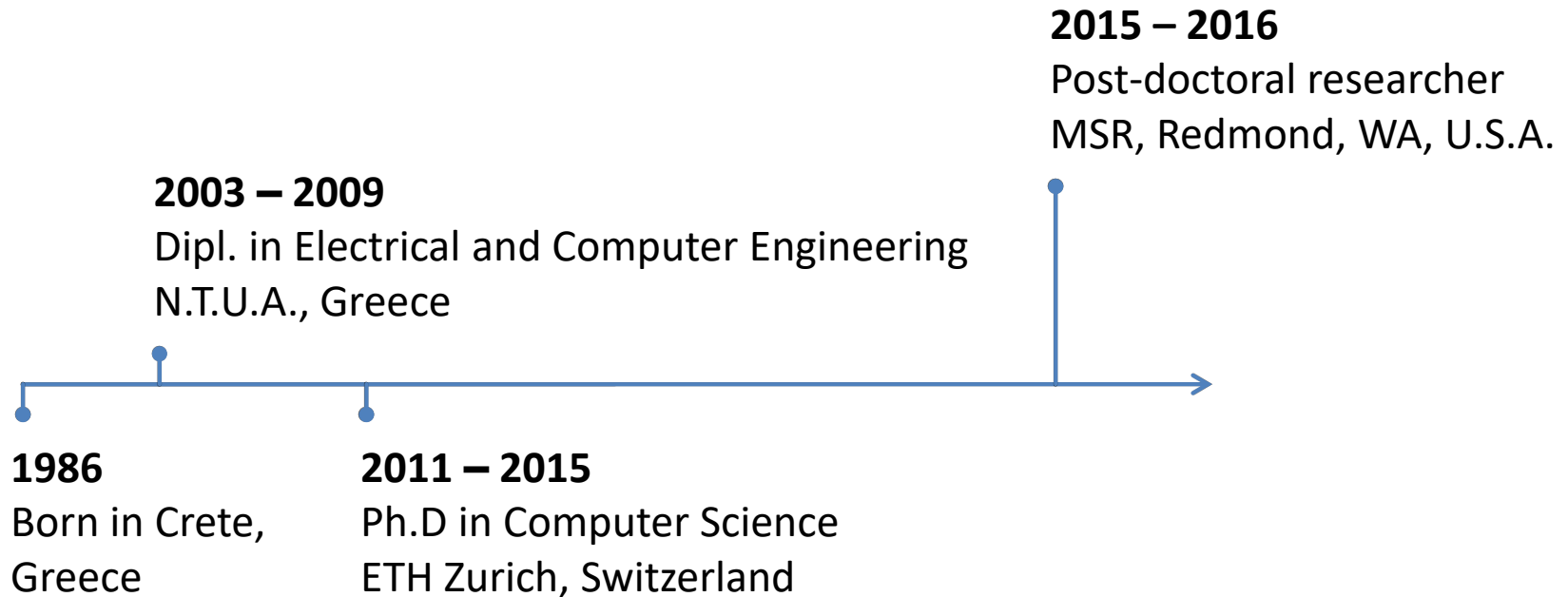
1986

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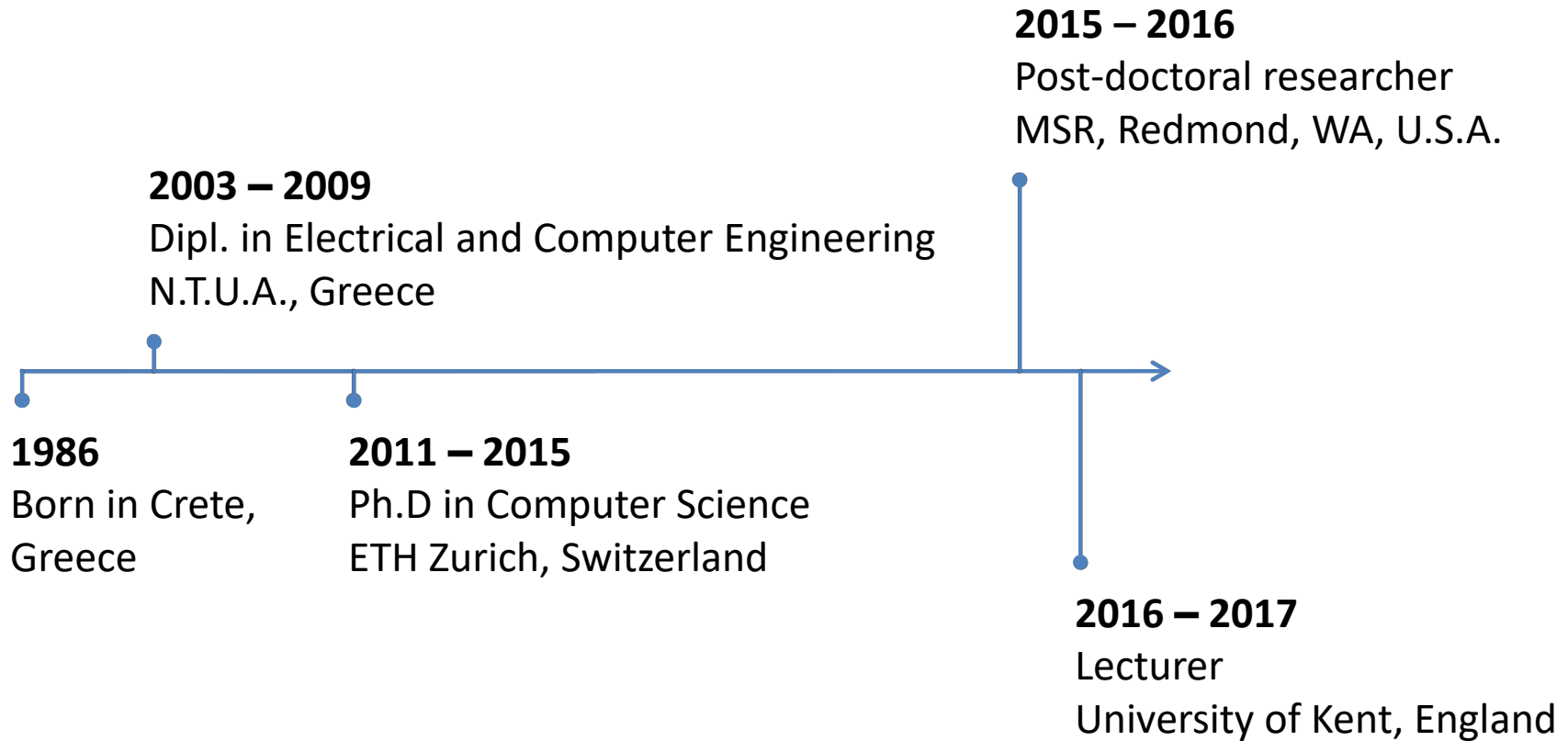
2011 – 2015

Ph.D in Computer Science
ETH Zurich, Switzerland

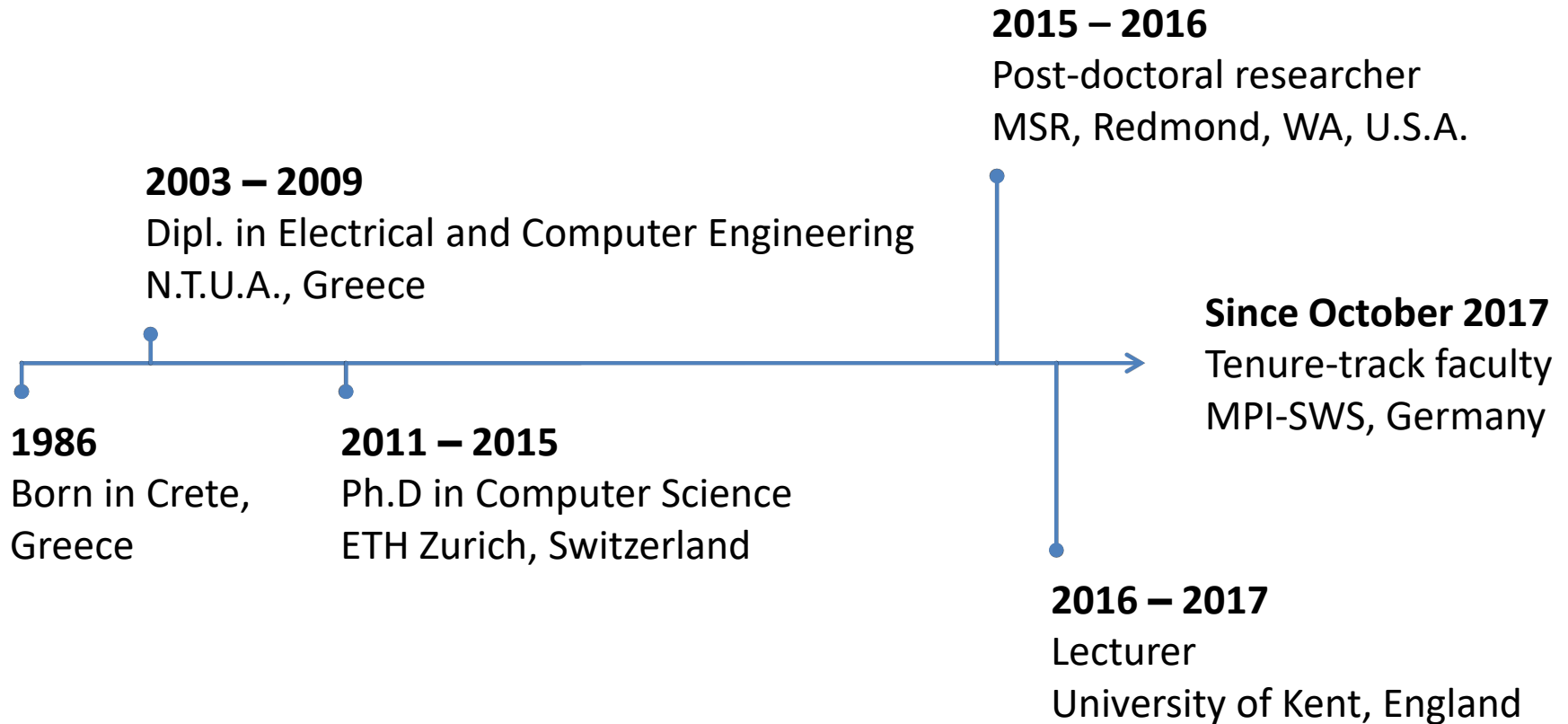
About me



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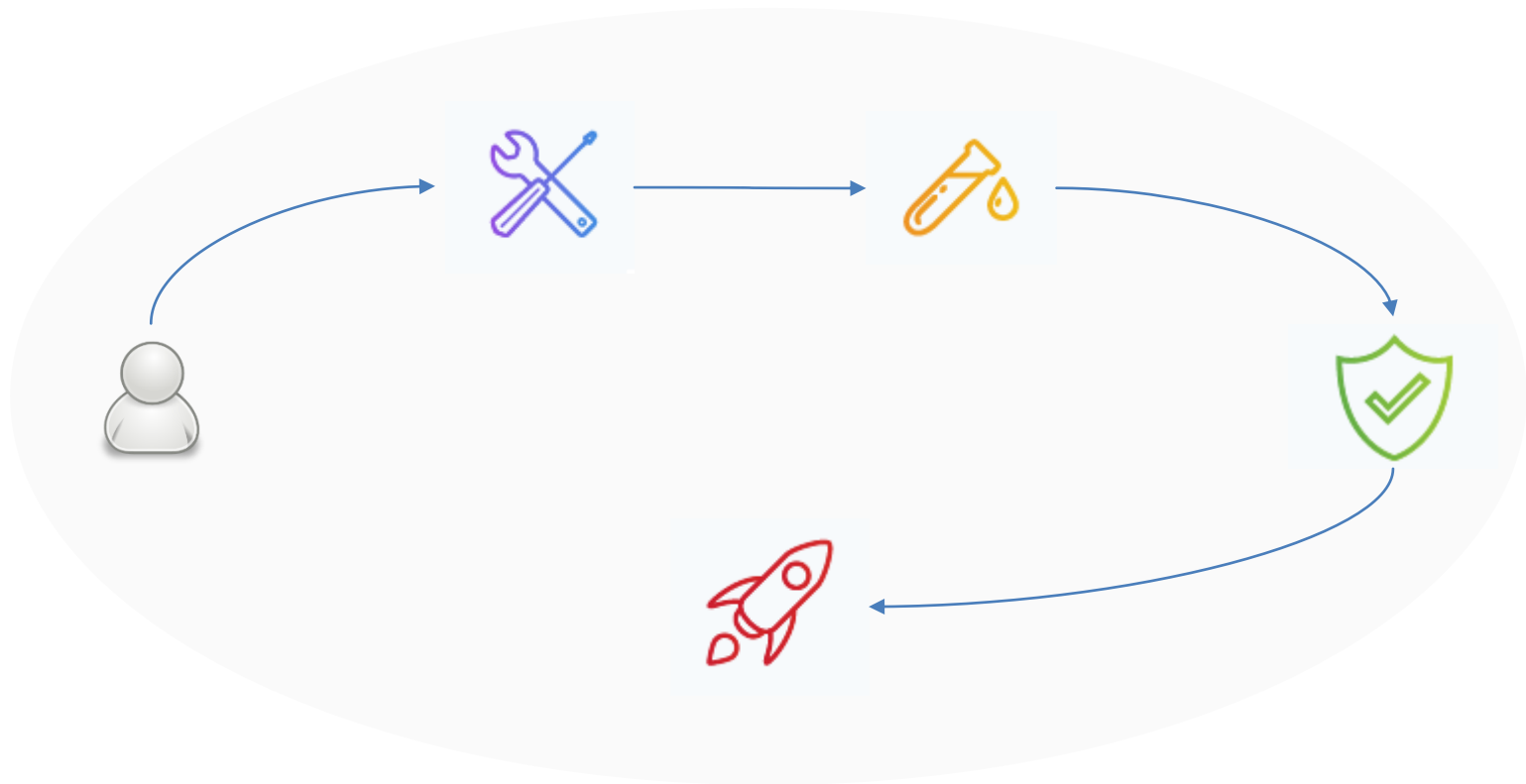


About me

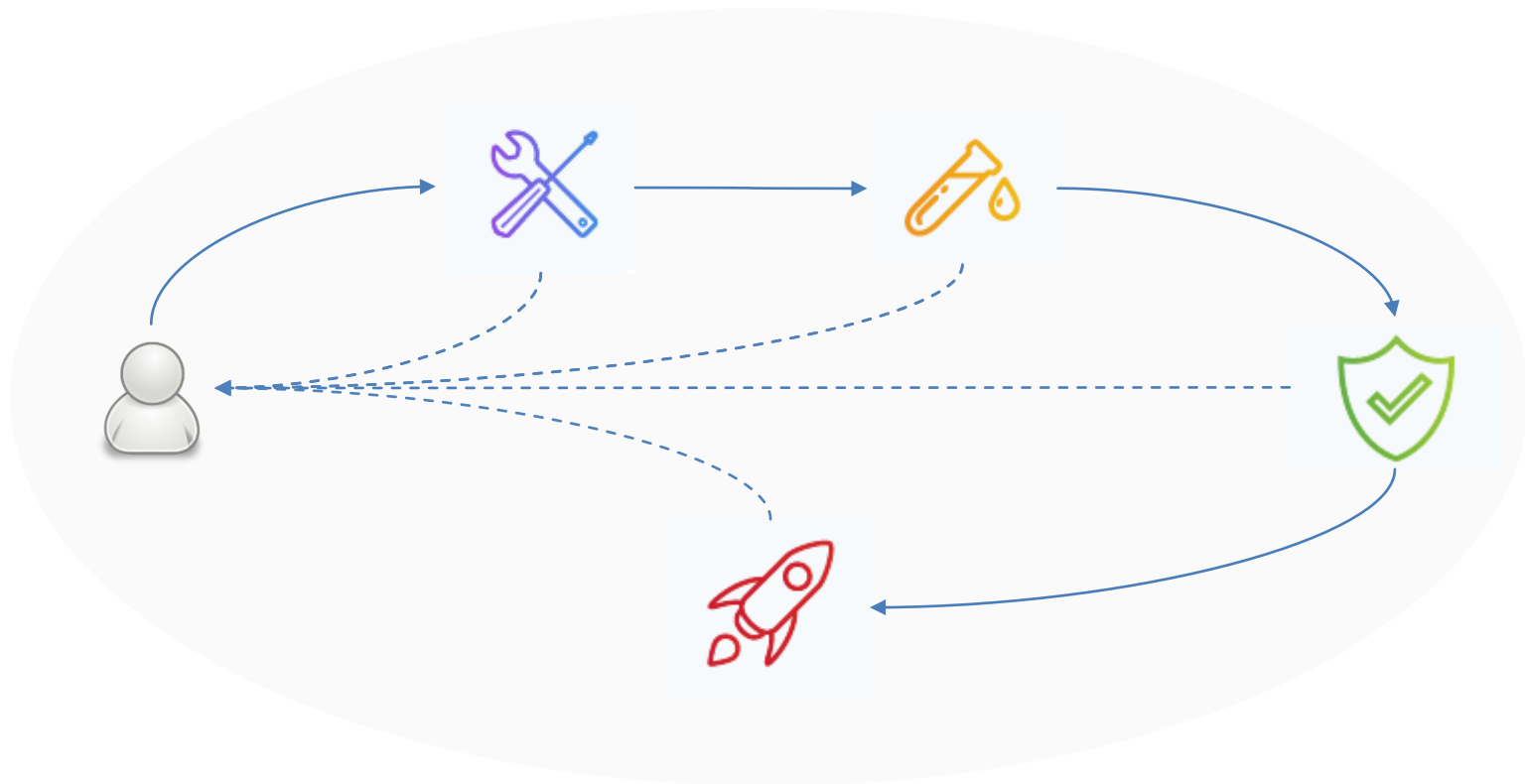


Practical formal methods

Practical formal methods



Practical formal methods



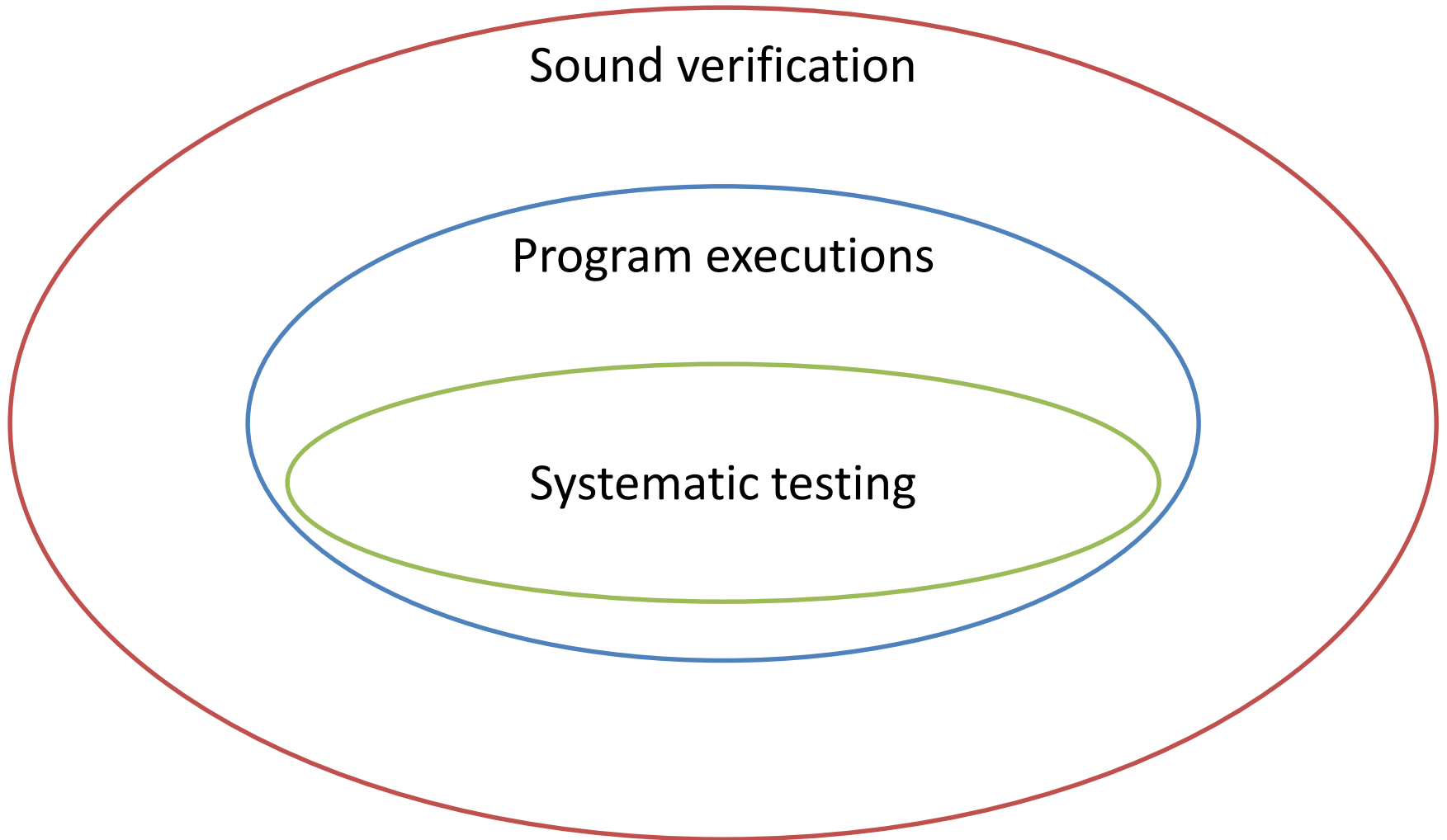
Our research mission

- Build more reliable software while increasing developer productivity
 - Formal foundation
 - Wide applicability
 - Strong practical impact

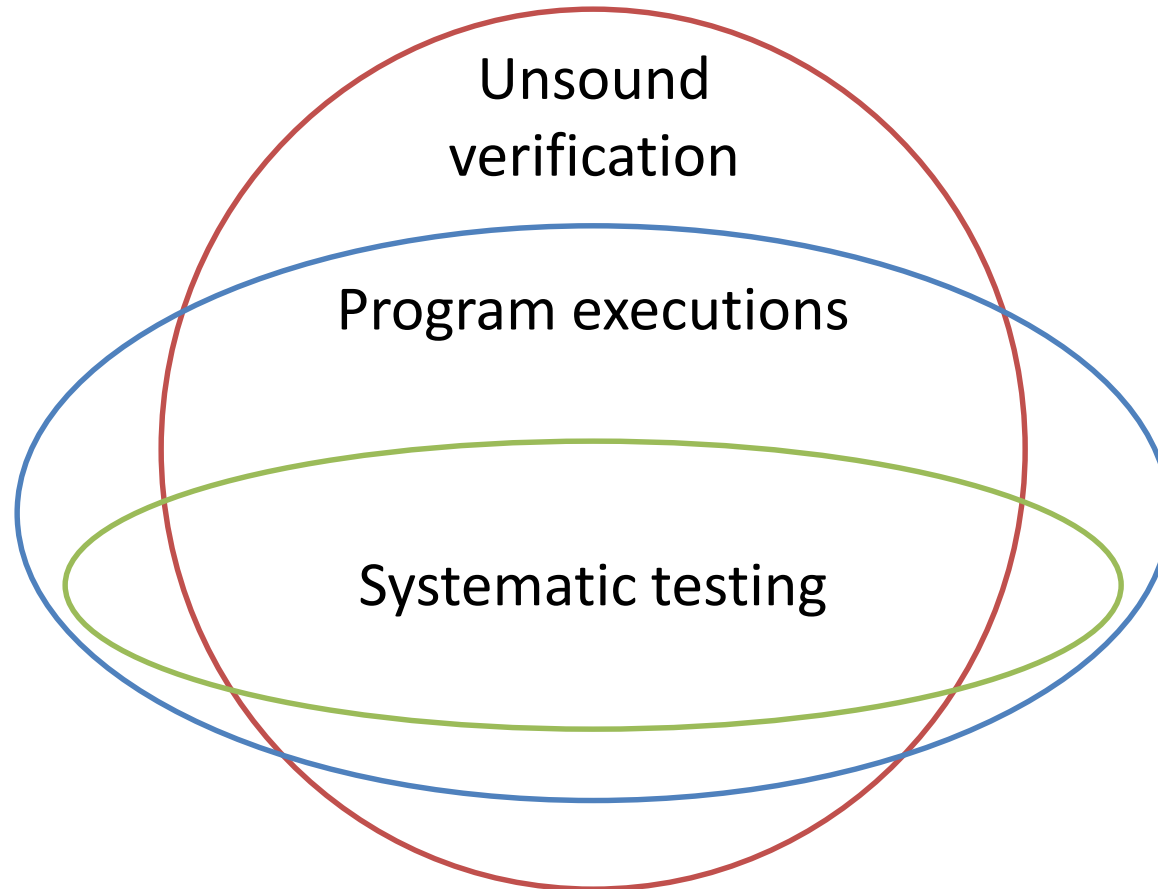
Our methodology

- Software engineering and formal methods
 - Automatic test generation
 - Program analysis and verification
 - Human-computer interaction
 - Empirical evaluation

Verification and testing

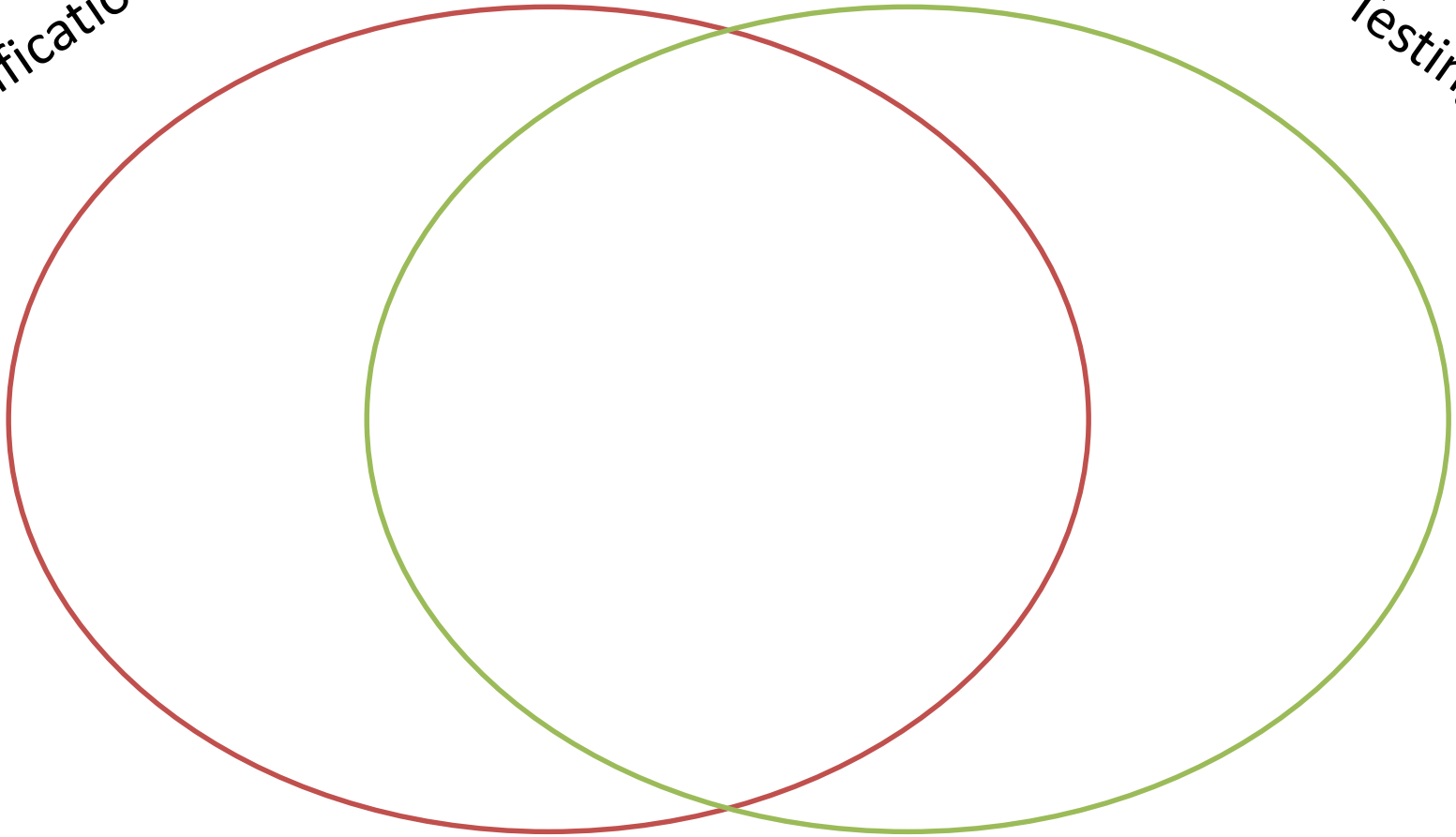


Verification and testing



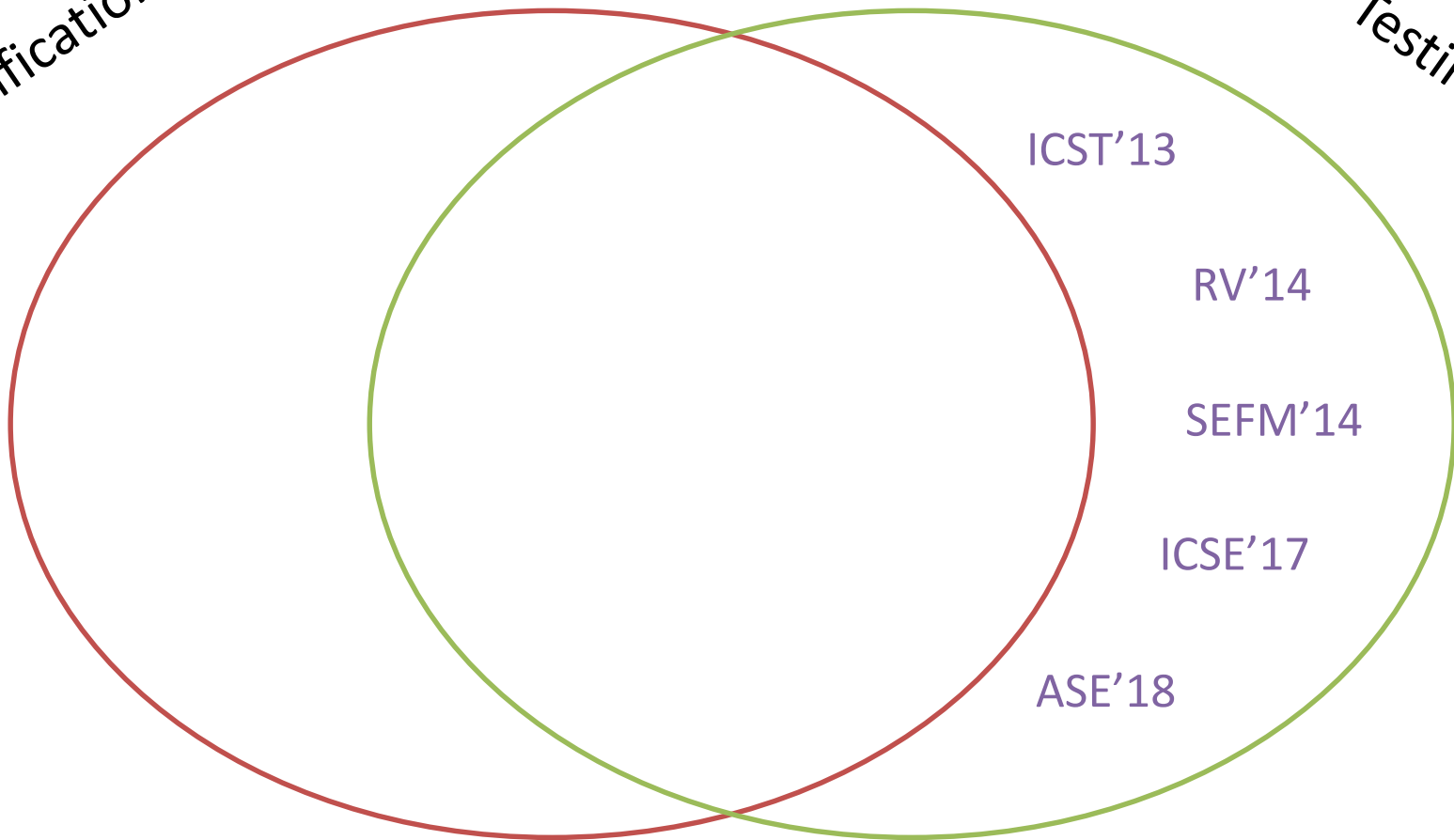
Verification

Testing



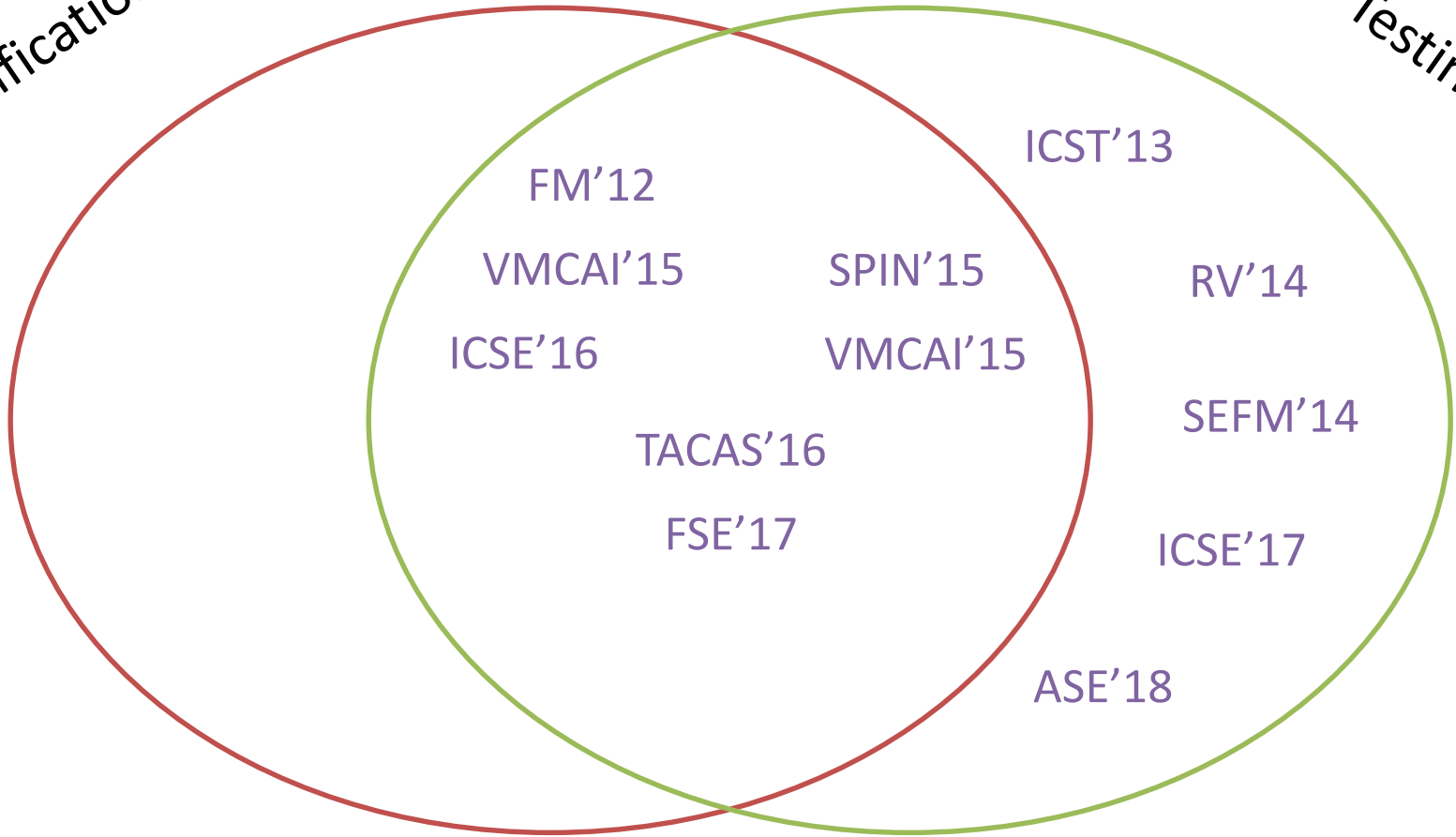
Verification

Testing



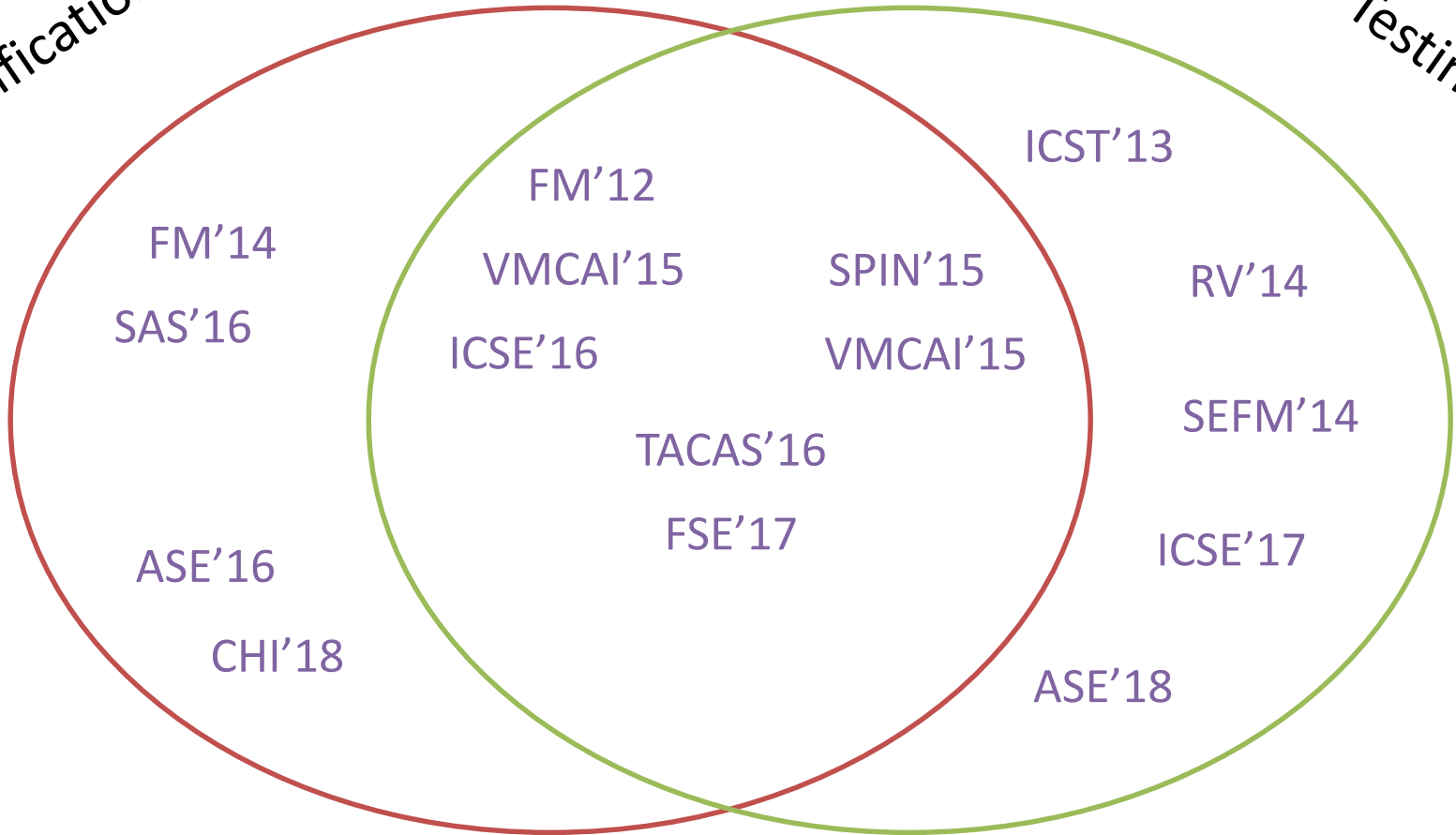
Verification

Testing



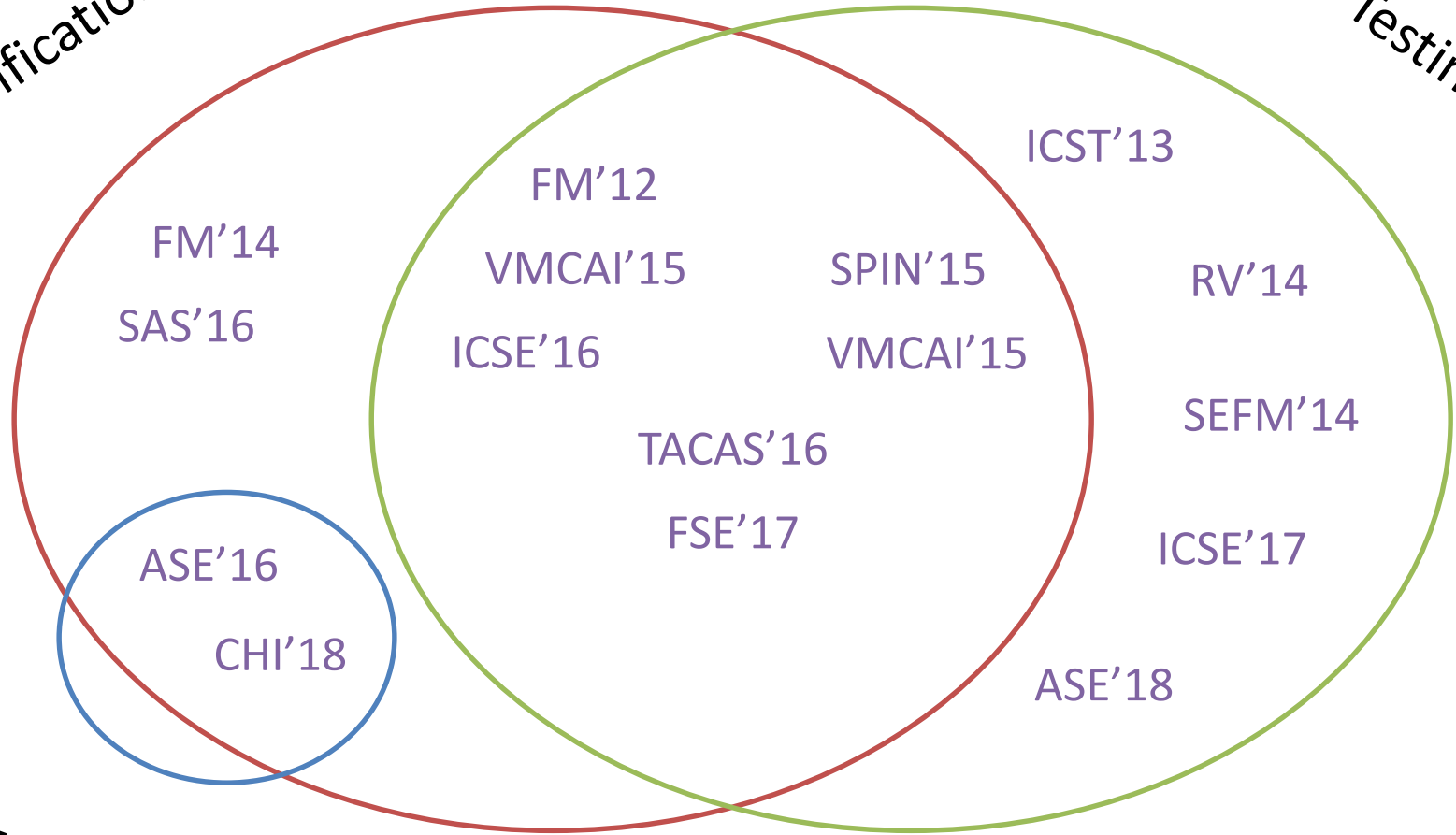
Verification

Testing



Verification

Testing



ESE & HCI

Static Program Analysis Meets *Test Case Generation*

Lecture 1

Textbook: Test Case Generation

Test case generation

- Detects bugs and security vulnerabilities
- Improves software quality
- Is cumbersome when done manually

Test case generation

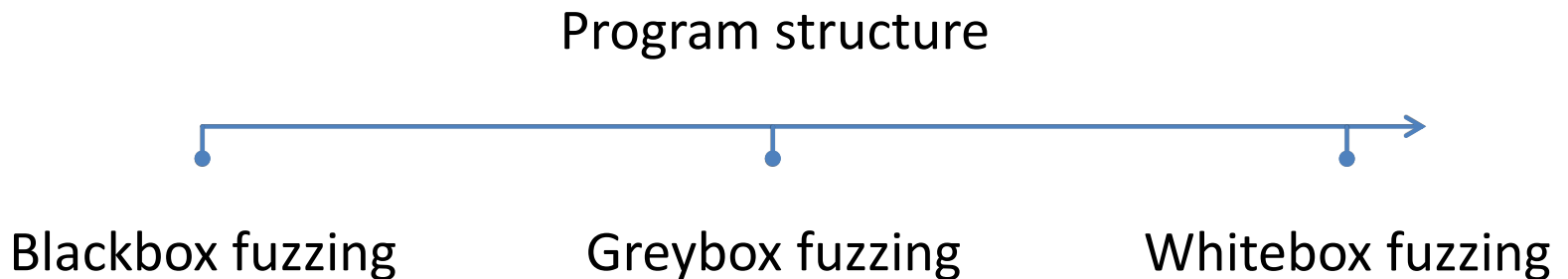
- Detects bugs and security vulnerabilities
- Improves software quality
- Is cumbersome when done manually

Characteristics

- Does not cover all program executions
- Does not generate false positives

Test case generation

- Detects bugs and security vulnerabilities
- Improves software quality
- Is cumbersome when done manually



Textbook: Blackbox Fuzzing

Blackbox fuzzing

- Also called random testing

Blackbox fuzzing

- Also called random testing

Algorithm

- Generates random inputs to a program
- Runs the program with these inputs
- Checks whether any bugs were found

Blackbox fuzzing: Discussion

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- Practical
- Very efficient

Blackbox fuzzing: Discussion

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but...

- Very ineffective in exploring new paths

Exercise

What is the probability of a blackbox fuzzer detecting the bug?

```
void foo(int x) {  
    if x == 42 { bug(); }  
}
```

Exercise

What is the probability of a blackbox fuzzer detecting the bug?

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void foo(int x) {  
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2^{32} !

Fun:

A General Framework for Dynamic Stub Injection

Background: Stub testing

- A technique to simulate dependencies of an application

Background: Stub testing

- A technique to **simulate dependencies** of an application
 1. Decide which **functions** to replace by stubs
 2. For each such function, implement one or more **stubs**
 3. For each **call** to such functions, decide whether to call a stub
 4. For each chosen call, perform the **stub injection**

Our framework

- Provides a domain specific language
 - Arbitrary stub behavior and flexible injection strategies
 - Stub reuse across programming languages
- Dynamically alters any executable program
 - No need for source code
- Detected 176 crashes in 10 commercial applications

A stub injection rule

```
kernel32.dll!GetModuleFileNameW(  
    hModule, lpFileName, nSize)
```

A stub injection rule

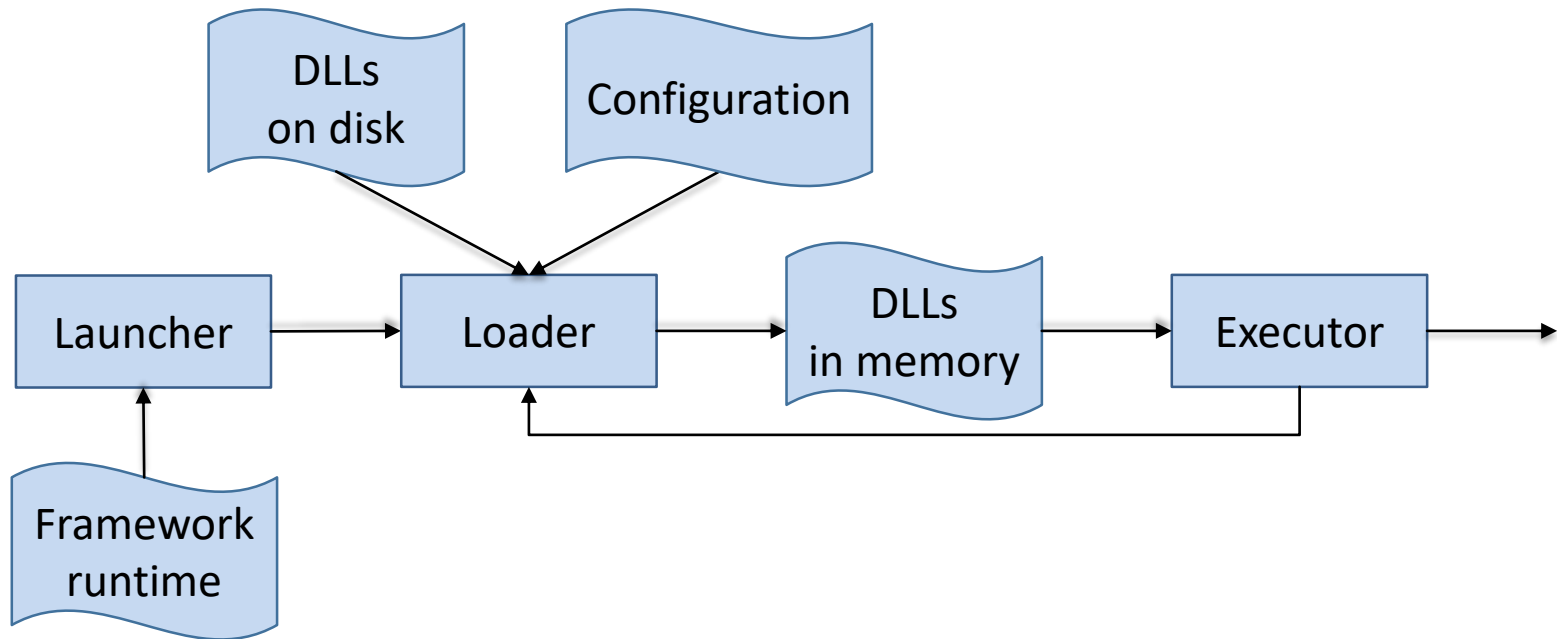
```
rule kernel32.dll!GetModuleFileNameW(  
    hModule, lpFileName, nSize)  
frequency every(2);  
after {  
    if (nSize < 32767) {  
        SetLastError(ERROR_INSUFFICIENT_BUFFER);  
        result = nSize;  
    }  
}
```

A stub injection rule

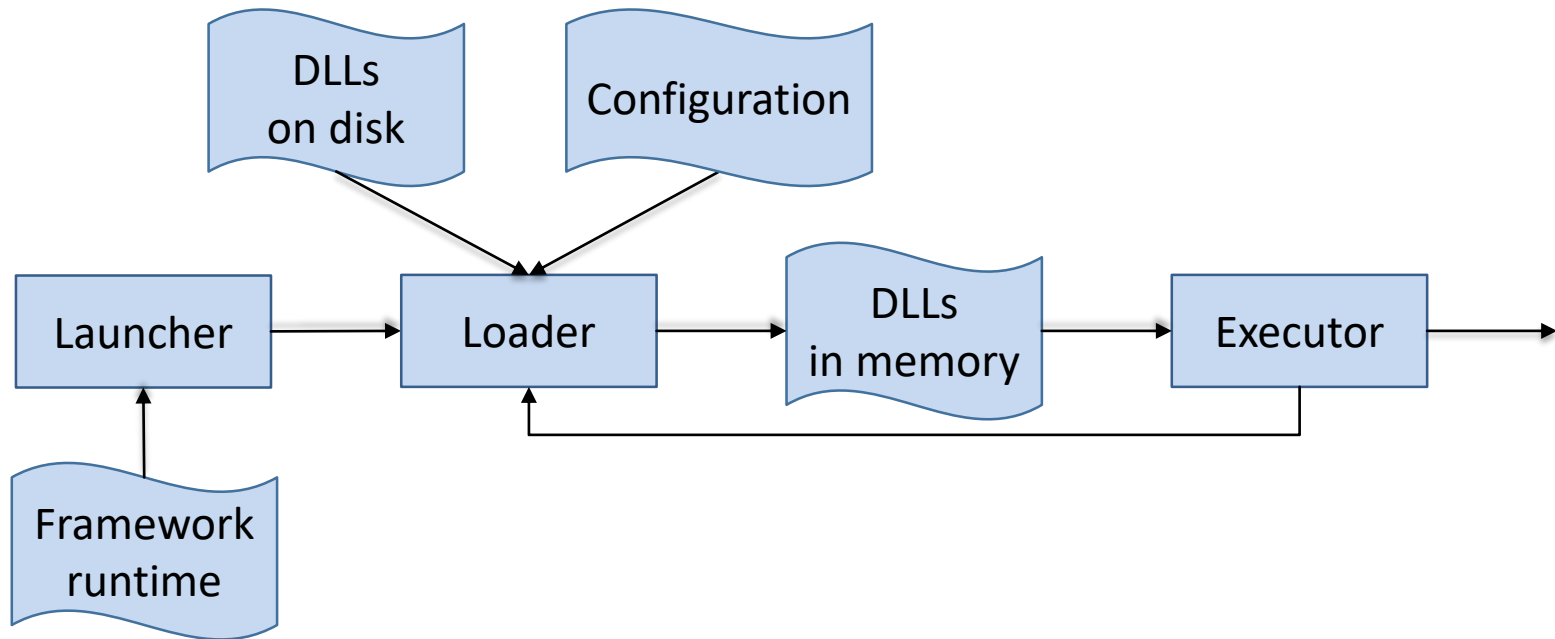
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frequency every(2);  
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}
```

- A stub can have non-trivial behavior
- A rule is language independent
- A stub can replace or augment the function behavior

The stub injection workflow



The stub injection workflow




- The instrumentation happens on the deployed code
- No need to know in advance which DLLs will be loaded
- Complex invocation schemes are easily handled

Experiment design: Fault injection

- Applications under test
- Test scenarios
- Stub injection rules

Experiment design: Fault injection

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- Test scenarios
- Stub injection rules

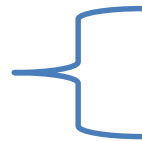


Internet Explorer Microsoft Excel
Microsoft Notepad
Microsoft OneNote Launcher
Microsoft Outlook Microsoft Paint
Microsoft PowerPoint
Microsoft Publisher
Microsoft Visio Microsoft Word
Notepad++ Skype for Business

Experiment design: Fault injection

- Applications under test

- Test scenarios



Start → Wait → Close

- Stub injection rules

Design: Stub injection rules

- Target functions
 - Requesting new resources
 - Accessing external resources
- Injection strategy
- Stub behavior

Design: Stub injection rules

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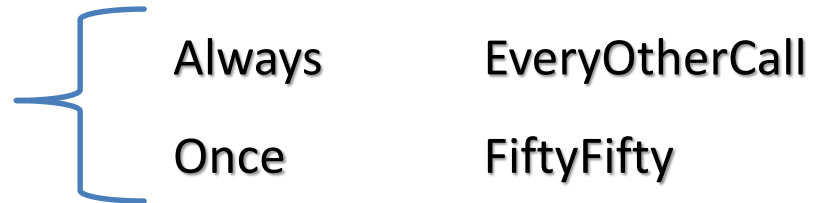
Microsoft Notepad

Notepad++

Design: Stub injection rules

- Target functions
 - Requesting new resources
 - Accessing external resources

- Injection strategy



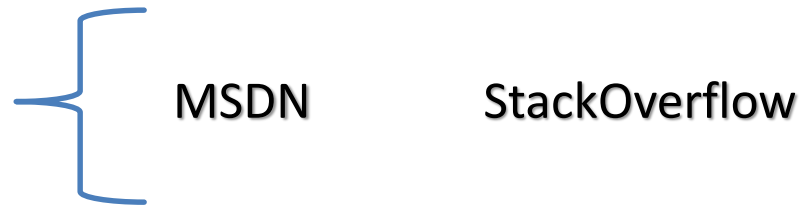
- Stub behavior

Design: Stub injection rules

- Target functions
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 - Accessing external resources

- Injection strategy

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Result highlights

- 160 rules (40 target functions x 4 injection strategies)
- 1,920 runs (160 rules x 12 applications)

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- 160 rules (40 target functions x 4 injection strategies)
- 1,920 runs (160 rules x 12 applications)
- 176 crashes in 10 applications
- 390,571 instrumented calls
- 140,621 stubbed calls

Results: Injection strategies

- All injection strategies detected crashes
- Always detects the most crashes, and Once the least
- Always may inject fewer stubs than EveryOtherCall and FiftyFifty
- Once detects by far the most crashes per 1,000 stubbed calls

Results: From crashes to bugs

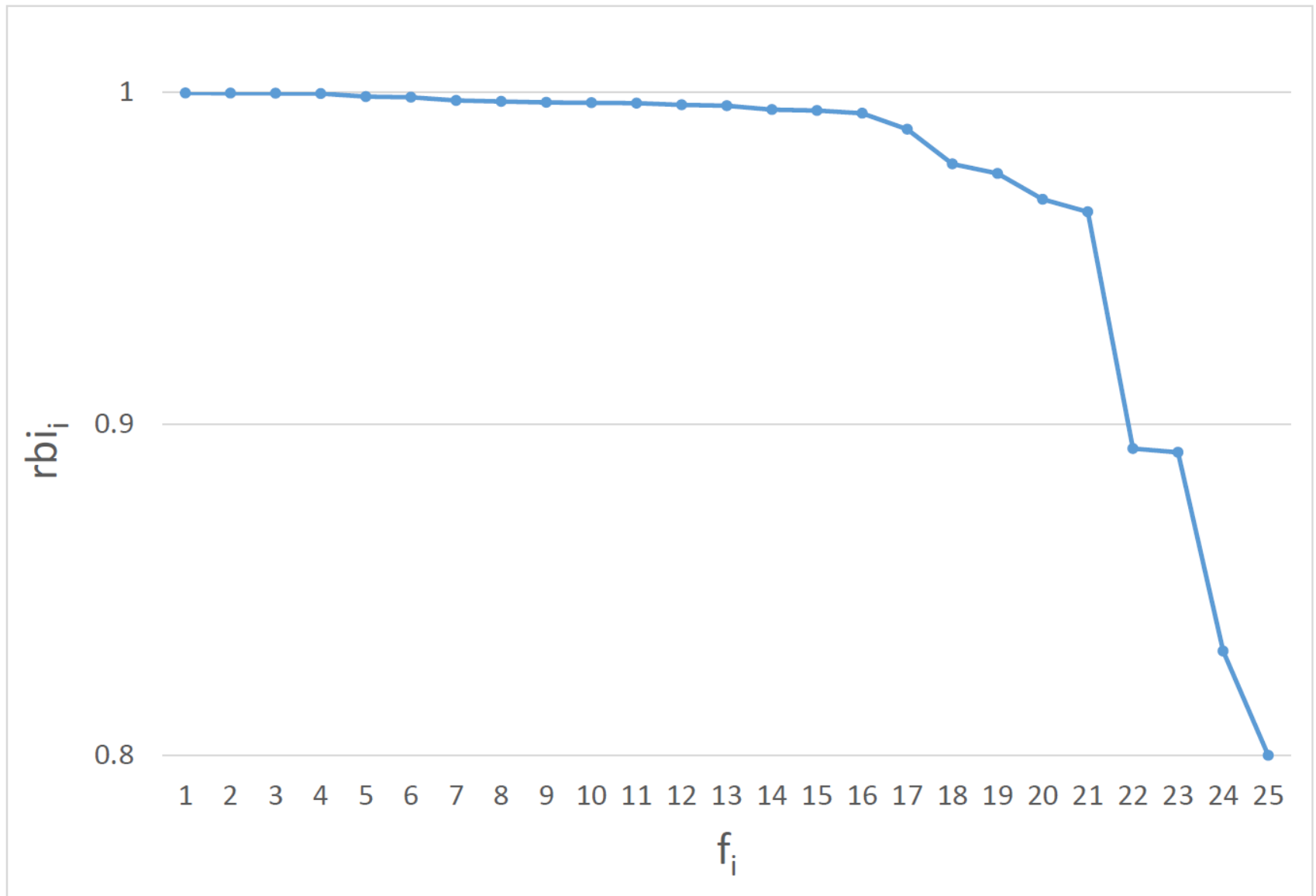
$$rbi_i = 1 - \frac{\text{number of crashes by stubbing } f_i}{\text{number of stubbed calls to } f_i}$$

Results: From crashes to bugs

$$rbi_i = 1 - \frac{\text{number of crashes by stubbing } f_i}{\text{number of stubbed calls to } f_i}$$

- Rules with a low rbi_i should be revised
- Developers should prioritize crashes by rules with high rbi_i
- New applications should be tested with rules with high rbi_i

Results: The real-bug indicator



Results: Bug in Excel

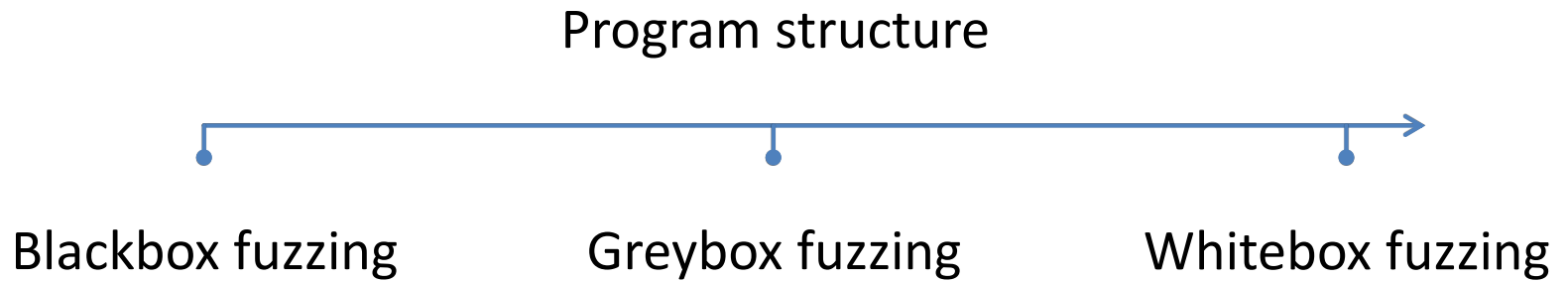
```
for (int c = 0; c < 2; c++) {
    error = RegQueryValueExW(key, /* ... */,
                             buffer, &buffer_size);
    if (ERROR_SUCCESS == error) {
        value_ptr = buffer;
        break;
    }
    else if (ERROR_MORE_DATA == error) {
        buffer.resize(buffer_size)
    }
}
assert(error == ERROR_SUCCESS);
```

Results: Rule for bug in Excel

```
rule ADVAPI32.dll!RegQueryValueExW(hKey,
    /* ... */, lpData, lpcbData)
before {
    if (last_key != hKey) {
        last_key = hKey;
        counter = 0;
    } else {
        counter++;
    }
}
after {
    if (counter < 2) {
        (*lpData)++;
        SetLastError(ERROR_MORE_DATA);
        result = ERROR_MORE_DATA;
    }
}
```

Textbook: Greybox Fuzzing


Greybox fuzzing



Greybox fuzzing: Algorithm

- Runs the program with an input
- Computes the identifier of the explored path
- Mutates the input and runs the program again
- Adds the new input to the test suite, if a new path was explored
- Picks a new input from the test suite

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Greybox fuzzing: Example

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int bar(int a, int b, int c) {
    var d = b + c;
    if d < 1 {
        if b < 3 {
            return 1;
        }
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a = -1, b = 0, c = -5

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a = -1, b = 0, c = -5
a = 0, b = 0, c = -5

Greybox fuzzing: Example

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a = -1, b = 0, c = -5
~~a = 0, b = 0, c = 5~~

Greybox fuzzing: Example

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int bar(int a, int b, int c) {  
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}
```

a = -1, b = 0, c = -5
~~a = 0, b = 0, c = 5~~
a = -1, b = 7, c = -5

Greybox fuzzing: Example

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int bar(int a, int b, int c) {  
    var d = b + c;  
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a = -1, b = 0, c = -5
~~a = 0, b = 0, c = 5~~
a = -1, b = 7, c = -5
a = -1, b = 7, c = 64

Greybox fuzzing: Example

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4 out of 5 paths
covered in 12 hours!

Greybox fuzzing: Discussion

Greybox fuzzing: Discussion

- Very practical
- Efficient

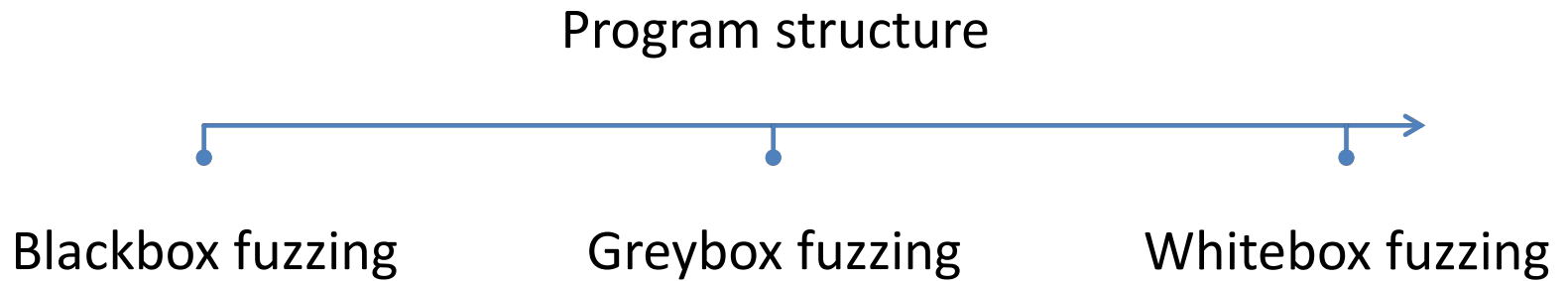
Greybox fuzzing: Discussion

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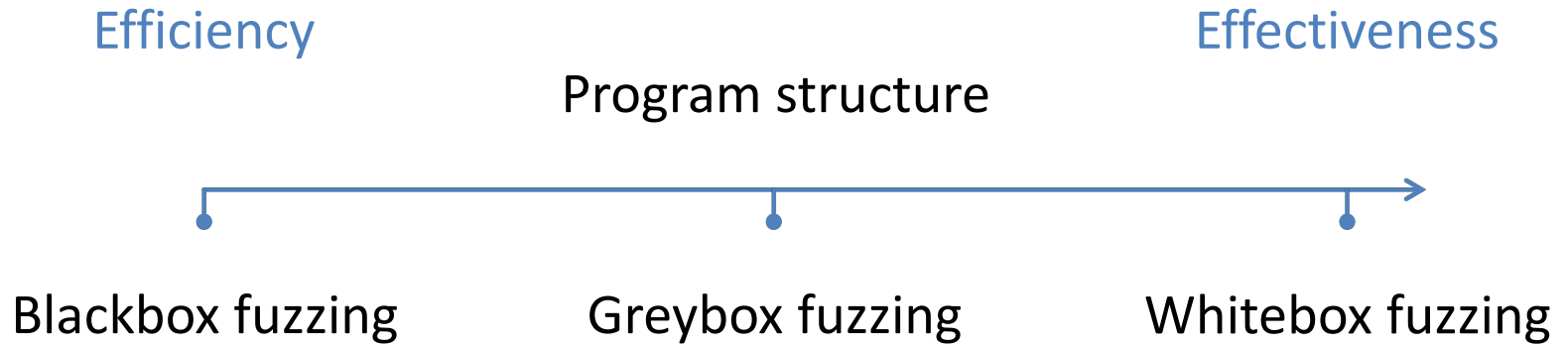
but...

- Ineffective in exploring new paths
(Inputs are still randomly mutated)

Greybox fuzzing: Discussion



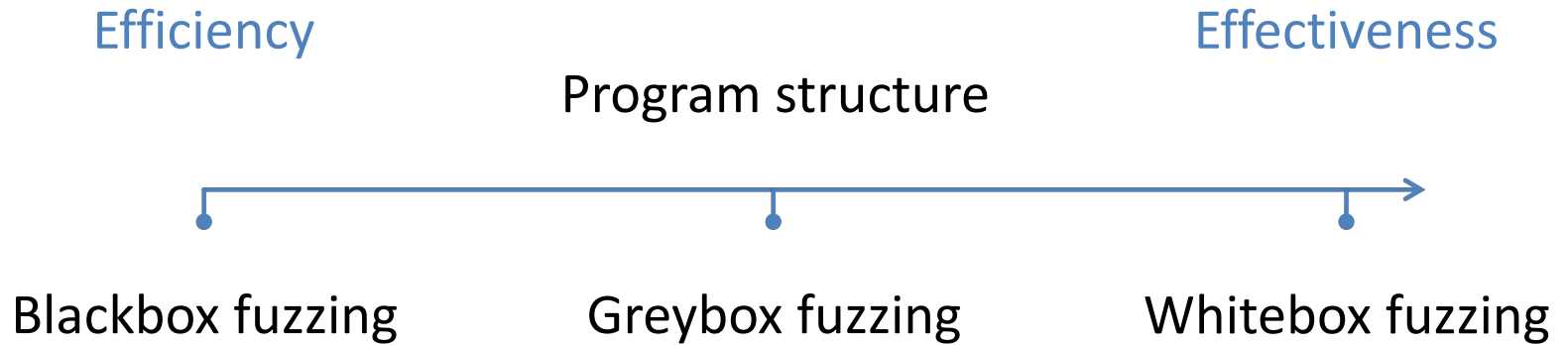
Greybox fuzzing: Discussion



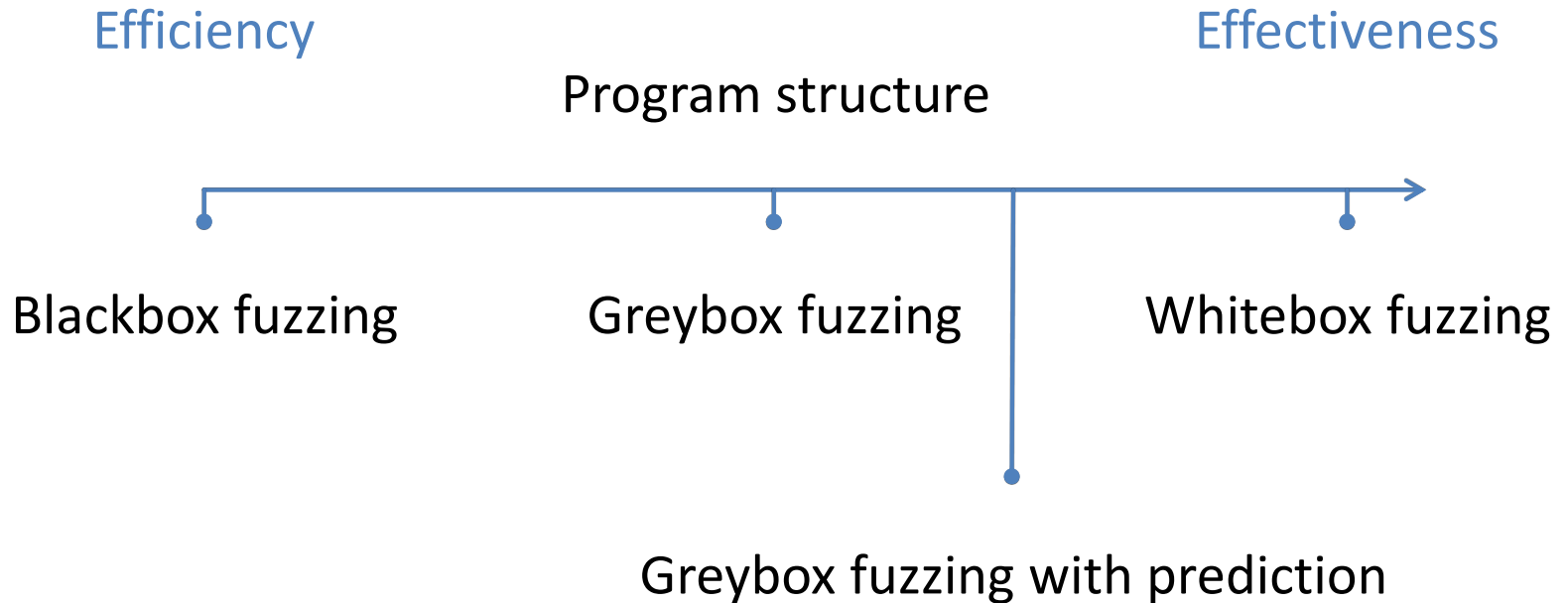
Fun:

Greybox Fuzzing with Cost-Directed Input Prediction

Greybox fuzzing



Greybox fuzzing: Input prediction

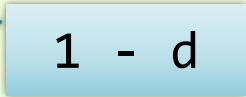


Input prediction: Example

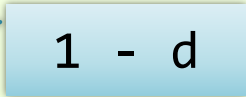

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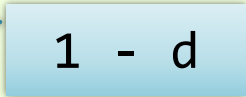


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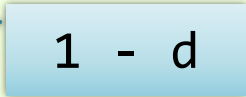



Input prediction: Example

```
int bar(int a, int b, int c) {  
    var d = b + c;  
    if d < 1 {  1 - d  
        if b < 3 {  
            return 1;  
        }  
        if a == 42 {  
            return 2;  
        }  
        return 3;  
    } else {  
        if c < 42 {  d  
            return 4;  
        }  
        return 5;  
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```

Input prediction: Example

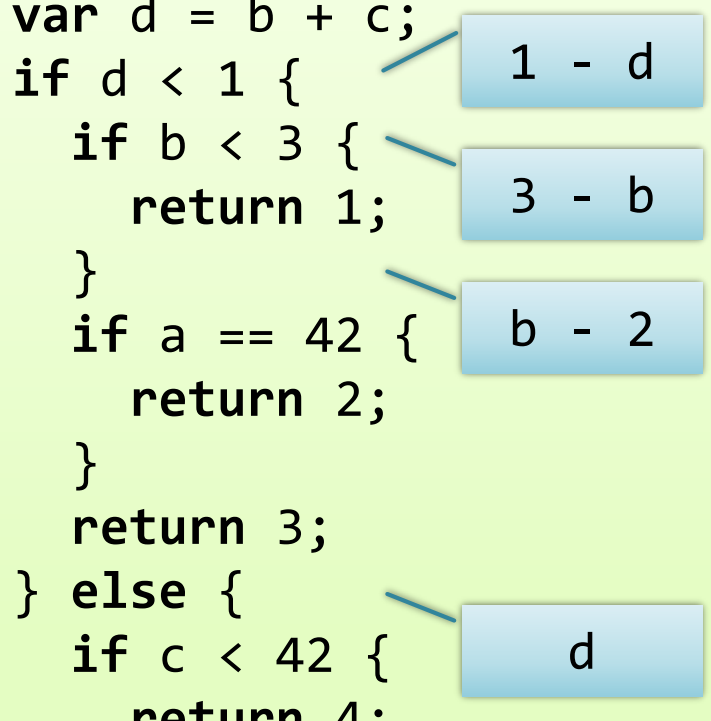
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```



1) $a = -1, b = 0, c = -5$

Input prediction: Example

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```

1 - d

3 - b

1) a = -1, b = 0, c = -5

Input prediction: Example

```
int bar(int a, int b, int c) {  
    var d = b + c;  
    if d < 1 { 1 - d c1:6  
        if b < 3 { 3 - b  
            return 1;  
        }  
        if a == 42 {  
            return 2;  
        }  
        return 3;  
    } else {  
        if c < 42 {  
            return 4;  
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        return 5;  
    }  
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```

1) $a = -1, b = 0, c = -5$

Input prediction: Example

```
int bar(int a, int b, int c) {  
    var d = b + c;  
    if d < 1 { 1 - d c1:6  
        if b < 3 { 3 - b c1:3  
            return 1;  
        }  
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```

1) $a = -1, b = 0, c = -5$

Input prediction: Example

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int bar(int a, int b, int c) {  
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- 1) $a = -1, b = 0, c = -5$
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Input prediction: Example

```
int bar(int a, int b, int c) {  
  var d = b + c;  
  if d < 1 { 1 - d c1:6 c2:9  
    if b < 3 { 3 - b c1:3  
      return 1;  
    }  
    if a == 42 {  
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    }  
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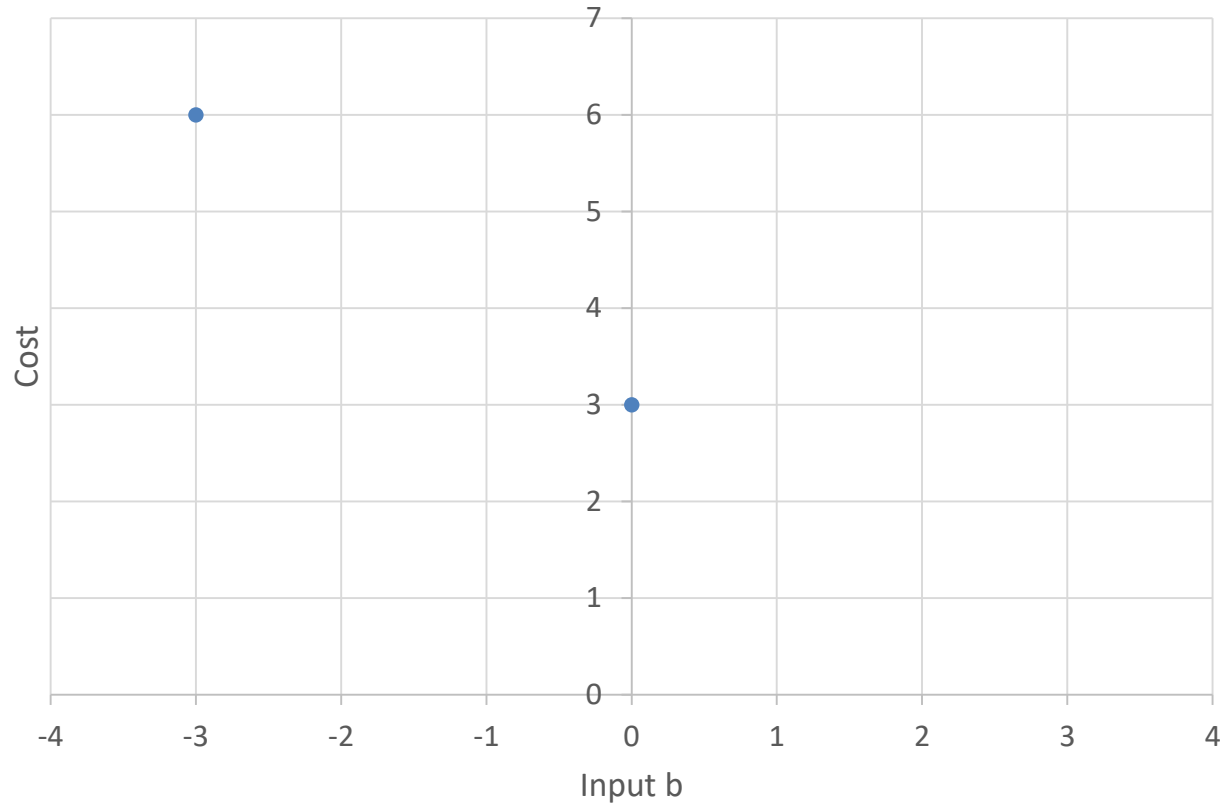
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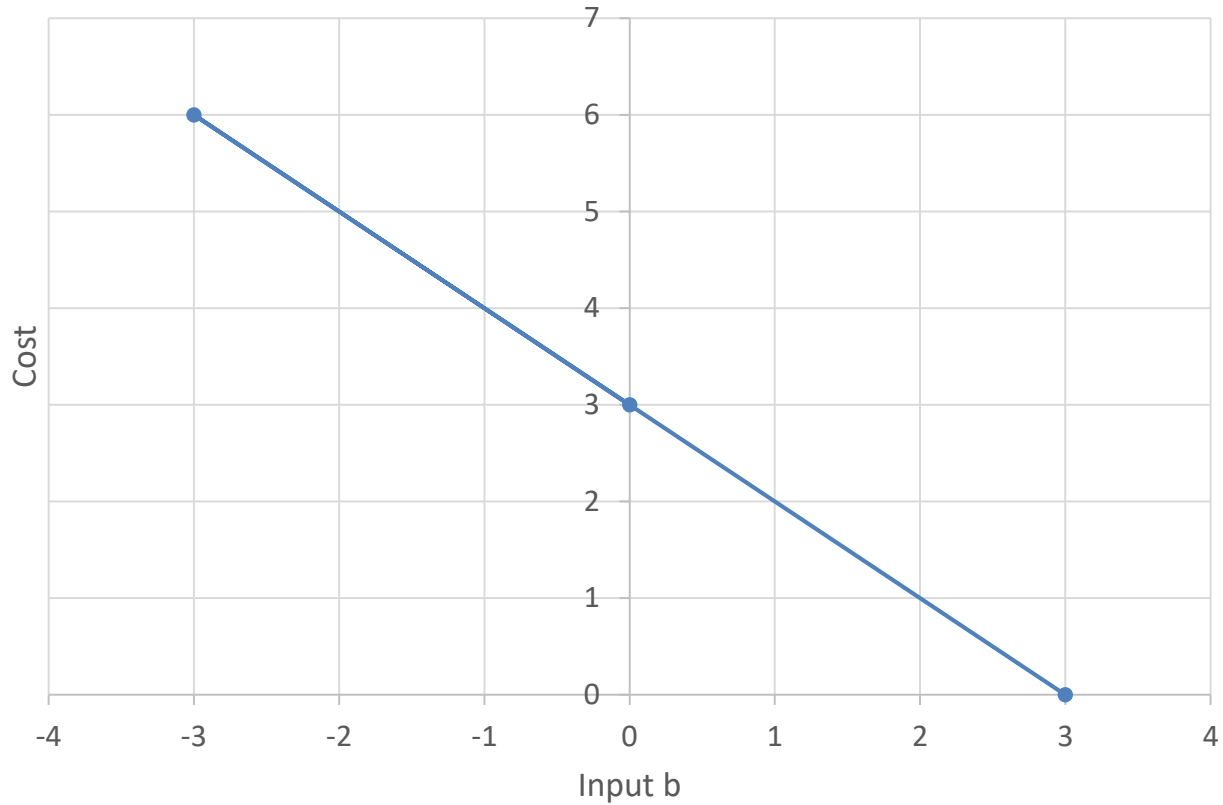
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Only 8 tests
for full path coverage!

Evaluation highlights

- Run on 26 applications, totaling 12'000 LOC
- Increased path coverage by up to 4X
- Found up to 77% more vulnerabilities, often orders-of-magnitude faster

Input prediction: Discussion

Input prediction: Discussion

- Very practical
- Efficient
- Effective in exploring new paths

Input prediction: Discussion

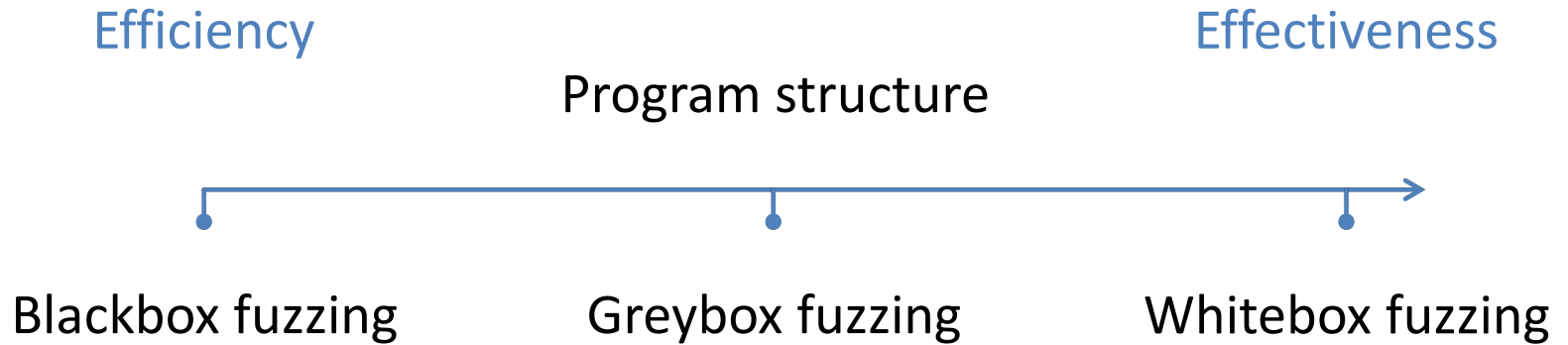
- Very practical
- Efficient
- Effective in exploring new paths

but...

- Approximate!

Textbook: Whitebox Fuzzing

Whitebox fuzzing



Whitebox fuzzing

- Also called **dynamic symbolic execution**

Whitebox fuzzing

- Also called **dynamic symbolic execution**

Algorithm

- Runs the program **concretely and symbolically**
- Collects constraints on inputs at branches
- Uses a constraint solver to generate new inputs

Whitebox fuzzing: Example

```
int bar(int a, int b, int c) {  
    var d = b + c;  
    if d < 1 {  
        if b < 3 {  
            return 1;  
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$a = -1, b = 0, c = -5$

$b + c < 1 \wedge b < 3$

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$a = -1, b = 0, c = -5$

$b + c < 1 \wedge b \geq 3$

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$a = -1, b = 0, c = -5$
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$b + c < 1 \wedge b \geq 3 \wedge \neg(a = 42)$

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$b + c \geq 1 \wedge c \geq 42$

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Whitebox fuzzing: Discussion

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- Practical
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but...

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Exercise

Encode the path constraints to determine the feasible paths.

```
int bar(int a, int b) {  
    var c = a * b;  
    if c >= 0 && b > 0 {  
        var d = c / 3 + a;  
        if d < 0 {  
            return 1;  
        }  
        return 2;  
    }  
    return 3;  
}
```

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        return 2;  
    }  
    return 3;  
}
```

2 feasible paths!

Exercise

How many paths are feasible?

```
void bar(int c) {  
    var i = 0;  
    while i < c {  
        i++;  
    }  
}
```


Exercise

How many paths are feasible?

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    }  
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```

2,147,483,647 + 1!

Fun:

Proving Memory Safety of the ANI Image Parser



Contributions

- First application of systematic testing for proving attacker memory safety
- of a complex and security-critical image parser
- using only compositional exhaustive testing
- How big is the gap between systematic testing and verification?

Attacker memory safety

- Means no attacker-controllable buffer overflows
- Is weaker than traditional memory safety
- Ignores input-independent memory accesses

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}
```

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untrusted
input

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void buggy(int x) {  
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    buf[x] = 1;  
}
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```
void buggy() {  
    char* buf = malloc(10);  
    buf[0] = 1;  
}
```

attacker memory unsafe

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trusted
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void buggy(int x) {  
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attacker memory unsafe

trusted
system call

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void buggy() {  
    char* buf = malloc(10);  
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```

attacker memory safe
but memory unsafe

SAGE

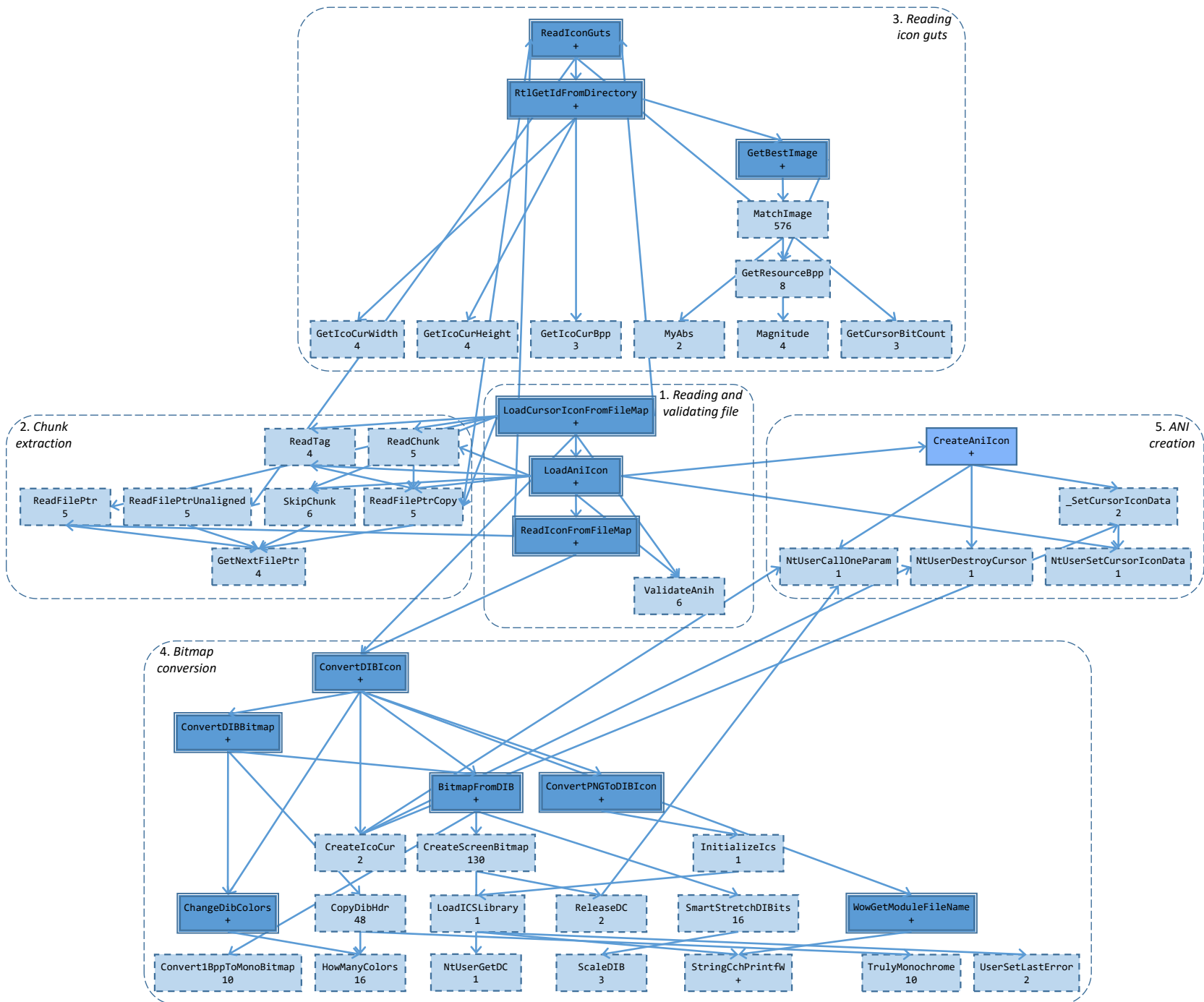
- Is a **whitebox fuzzer** for security testing
- Applies **dynamic symbolic execution** at the x86 level
- Is extensively used to fuzz several **Microsoft** products

Verification with testing?

- All program paths may be enumerated with **dynamic symbolic execution** when
 - there are finitely many paths
 - all feasible branches are explored
 - the constraint solver is sound and complete
- We turned maximum precision of SAGE **on** and any unsound pruning **off**

The ANI Windows image parser

- Implemented in C and x86 assembly
- At least 350 functions in 5 DLLs
- Testing untrusted inputs in at least 110 functions
- 47 top-level functions in `user32.dll`
- No recursion and no complex data structures



Strategies

- Stage 1: Bottom-up strategy
 - Exhaustive path enumeration and inlining
- Stage 2: Input-dependent loops
 - Restricting input-dependent loop bounds
- Stage 3: Top-down strategy
 - Summarization

Strategies

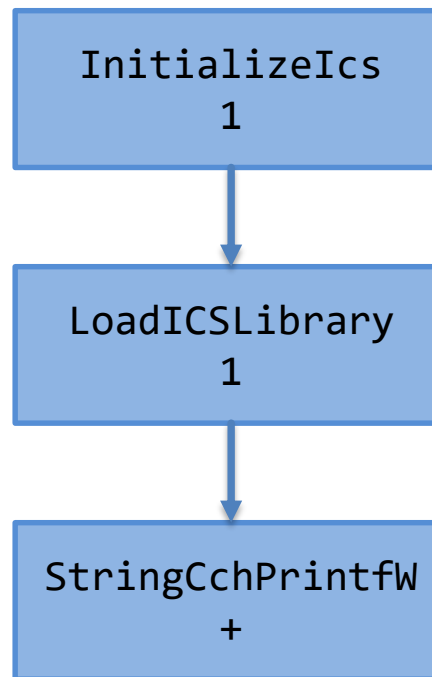
- Stage 1: Bottom-up strategy
 - Exhaustive path enumeration and inlining
- Stage 2: Input-dependent loops
 - Restricting input-dependent loop bounds
- Stage 3: Top-down strategy
 - Summarization
- Fixing or ignoring attacker-memory-safety bugs

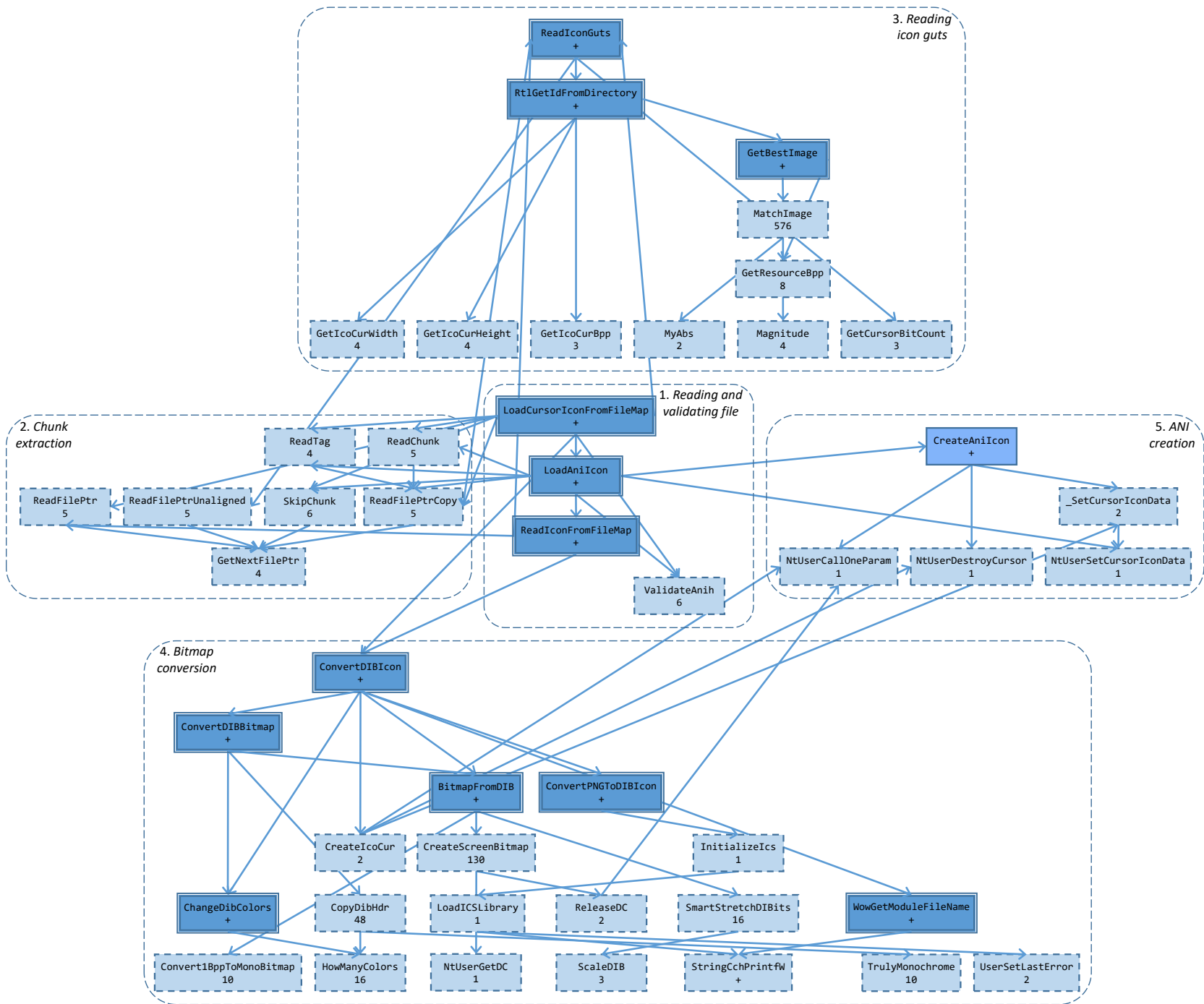
Stage 1: Bottom-up strategy

- Exhaustive path enumeration for all calling contexts in less than 12 hours

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34 out of 47 functions were verified!

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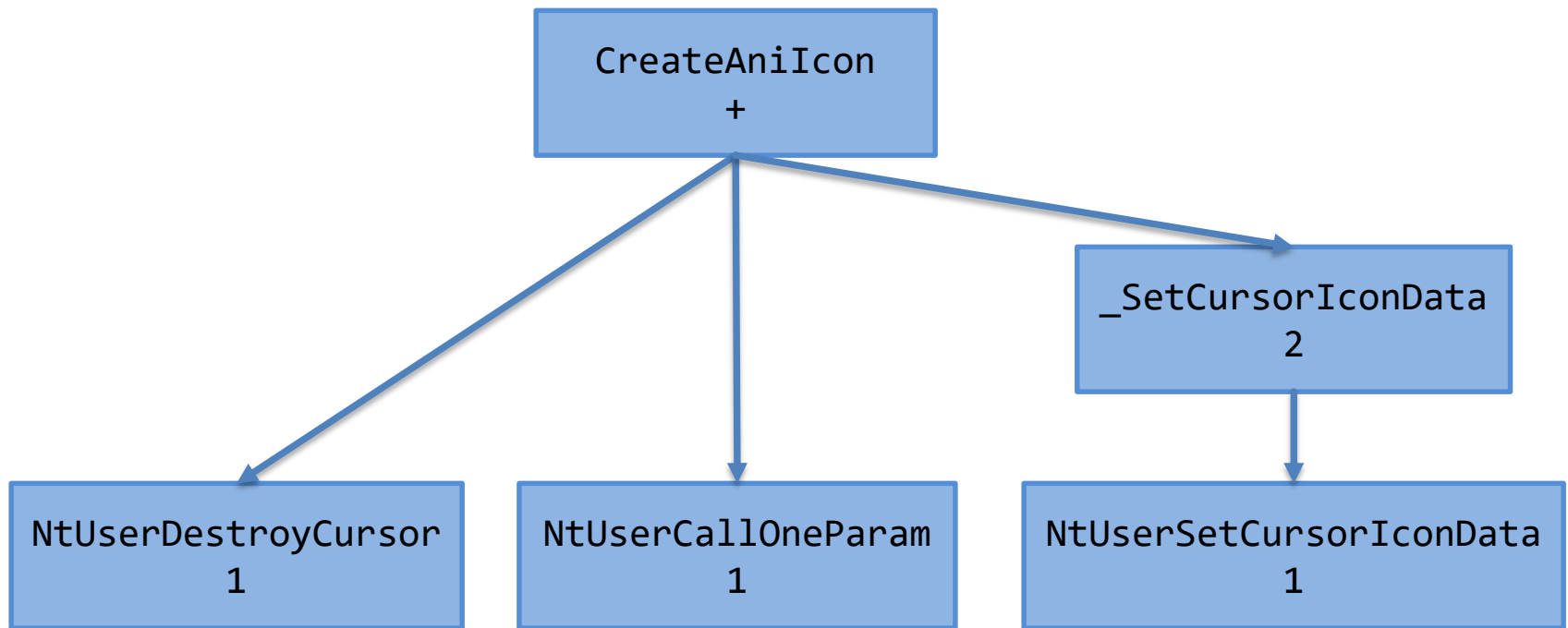


Stage 2: Input-dependent loops

- Manually restricting input-dependent loop bounds

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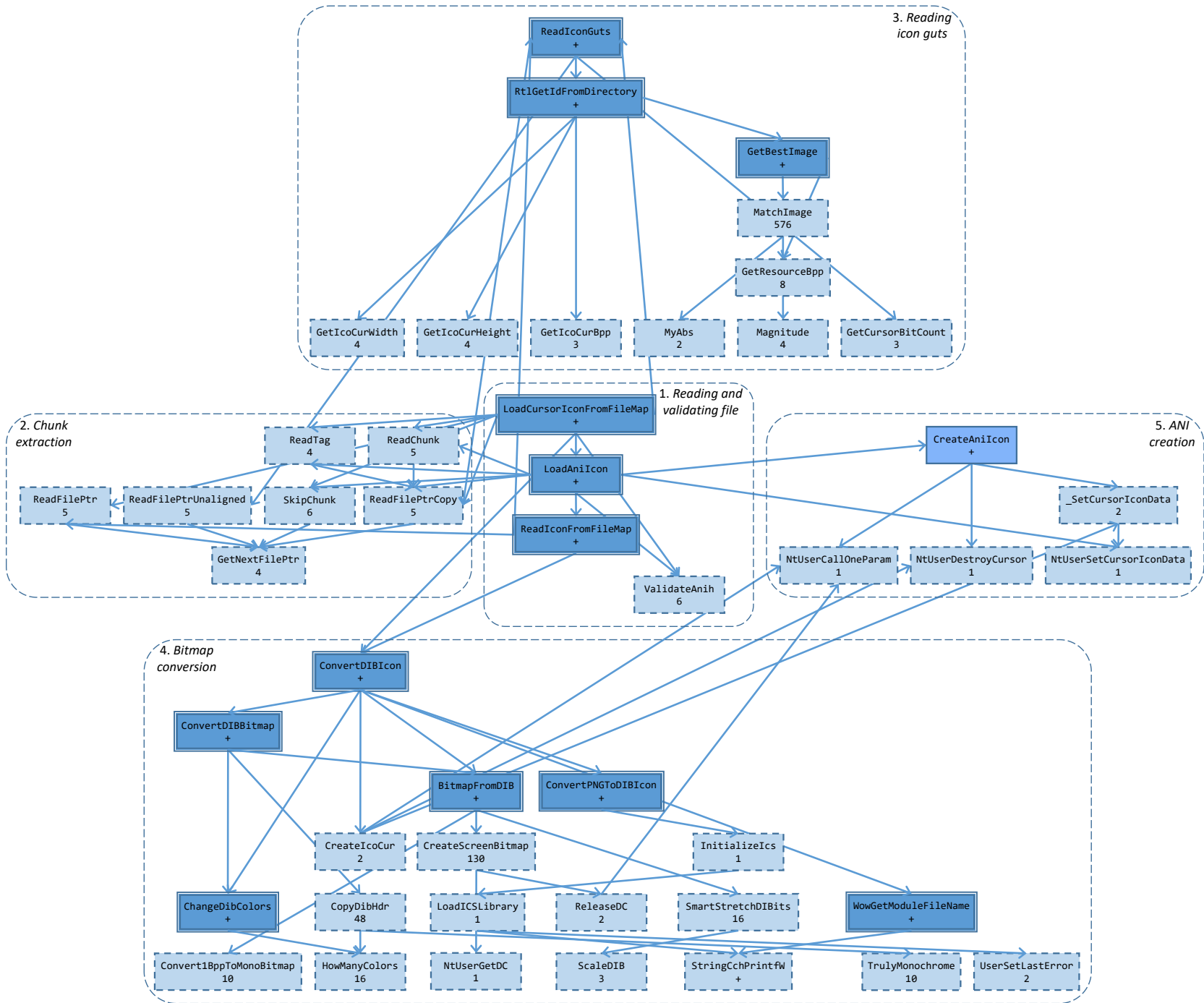
- Manually restricting input-dependent loop bounds



Stage 2: Input-dependent loops

- Manually restricting input-dependent loop bounds

```
for (i = 0; i < frames; i++)
    frameArrT[i] = frameArr[i];
for (i = 0; i < steps; i++) {
    if (rateArr == NULL)
        rateArrT[i] = rate;
    else
        rateArrT[i] = rateArr[i];
    if (stepArr == NULL)
        stepArrT[i] = i;
    else
        stepArrT[i] = stepArr[i];
}
```



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1 out of 47 functions was tested!

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10 input bytes plus the file size!

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Stage 3: Top-down strategy

- Manual program decomposition and summarization

Summarization

- Alleviates path explosion at the expense of computing logic representations of functions

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- Alleviates path explosion at the expense of computing logic representations of functions
- Is attractive only for functions
 - that contain many execution paths, and
 - whose inputs/outputs are easy to represent logically

Compositional memory safety

```
void bar(char* buf, int x) {  
    if (0 <= x && x < 10) buf[x] = 1;  
}
```

- Summary for bar:

$$\left((0 \leq x) \wedge (x < 10) \wedge (0 \leq x < \text{buf.size}) \wedge (\text{buf}[x]=1) \right) \vee$$
$$\left((x < 0) \wedge (10 \leq x) \right)$$

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void foo(int x) {  
    char* buf = malloc(5);  
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Compositional memory safety

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- Summary for bar:

5

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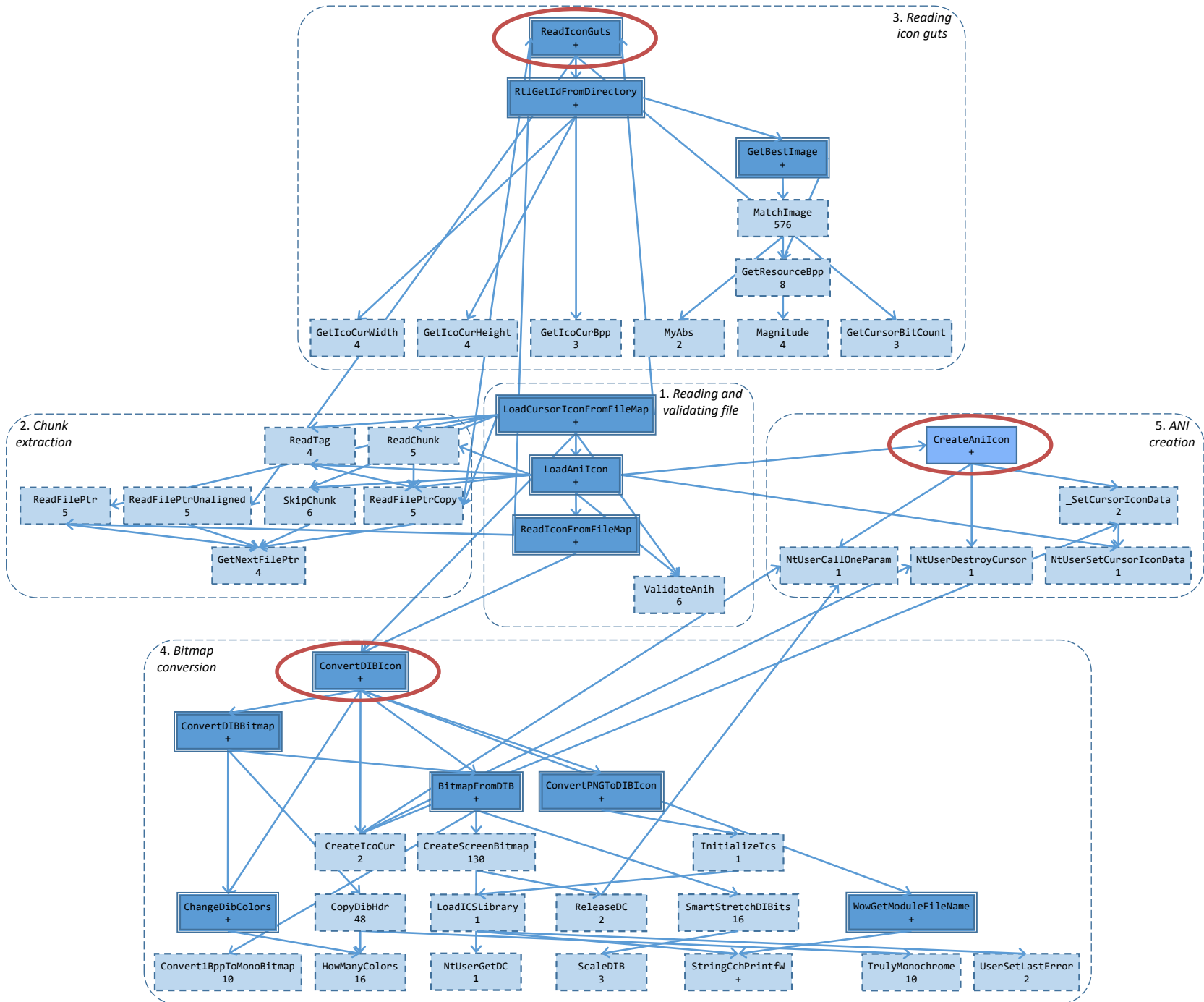
```
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    char* buf = malloc(5);  
    bar(buf, x);  
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```

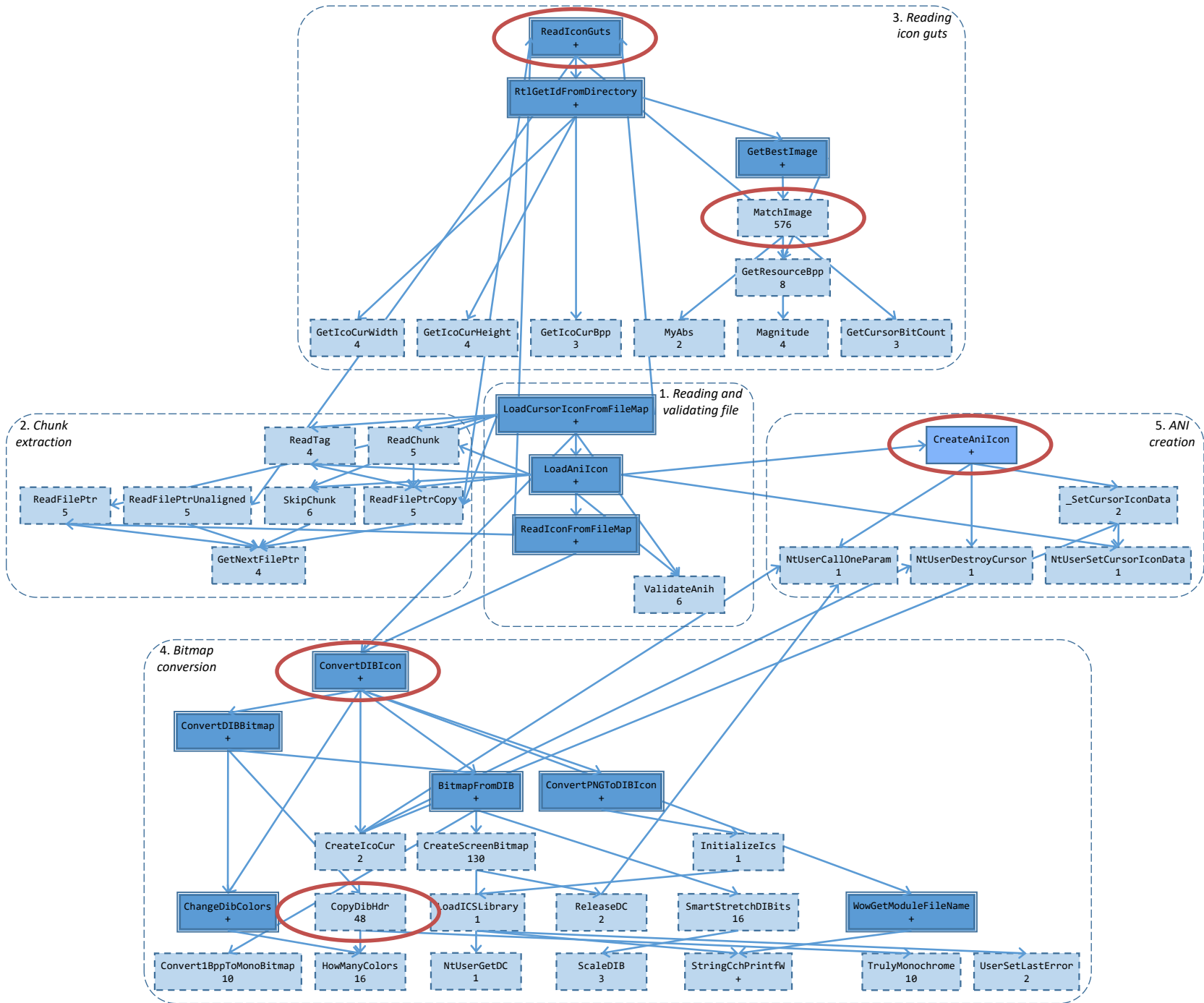
Stage 3: Top-down strategy

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Stage 3: Top-down strategy

- Manual program decomposition and summarization
- Key observation:
Architectural components of the parser typically return a single “success” or “failure” value
 - If “failure”, parsing aborts
 - If “success”, parsing proceeds with unconstrained inputs





Stage 3: Top-down strategy

- Manual program decomposition and summarization

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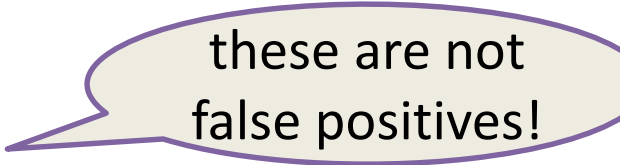
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Only 5 summaries were used!

Attacker-memory-safety bugs

- Real bugs (fixed in the latest version of Windows)
- Harmless (off-by-one) bugs
- Code parts not memory safe by design

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Shrinking the gap

- Interesting findings:
 - Many loop-free and easy-to-verify functions
 - Input-dependent loops controlled by about 10 bytes plus the file size
 - Path explosion controlled by 5 simple summaries
- Can we automate the manual steps?

Static Program Analysis Meets Test Case Generation

Lecture 1

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